

MEMORANDUM
RM-3578-PR
AUGUST 1964

ON DISTRIBUTED COMMUNICATIONS:
III. DETERMINATION OF PATH-LENGTHS IN
A DISTRIBUTED NETWORK

J. W. Smith

PREPARED FOR:
UNITED STATES AIR FORCE PROJECT RAND

The **RAND** *Corporation*
SANTA MONICA • CALIFORNIA

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PREFACE

This Memorandum is one in a series of eleven RAND Memoranda detailing the Distributed Adaptive Message Block Network, a proposed digital data communications system based on a distributed network concept, as presented in Vol. I in the series.* Various other items in the series deal with specific features of the concept, results of experimental modelings, engineering design considerations, and background and future implications.

The series, entitled On Distributed Communications, is a part of The RAND Corporation's continuing program of research under U.S. Air Force Project RAND, and is related to research in the field of command and control and in governmental and military planning and policy making.

The present Memorandum, the third in the series, is a continuation of the model simulation study reported in the previous volume. Since a network of the type proposed had never been built, there was much we did not know about its performance. For example, we wanted to learn more about the distribution of message path-lengths in the network, in order that transmission times might be determined; we wanted to know how the network behaved under heavy loading--such as would occur during a crisis;

*A list of all items in the series is found at the end of the Memorandum.

we wanted to know how the network would react when a "hog" station or stations attempted to purposely overload the network; we wanted to know how many message blocks would be lost if a policy of dropping traffic that has circulated longer than some specified time were used.

Because of the complexity of such networks, our only high-confidence tool for predicting performance was the Monte Carlo simulation, wherein messages are created and circulated in a computer model of the network, statistics of traffic-flow are examined, network parameters are changed, and the network is re-examined.

A FORTRAN simulation (for the IBM 7090 computer) was described in Vol. II in the series, together with the results of the simulation performed. There were three shortcomings to this effort: first, the size network that could be accommodated was smaller than desired; second, an undue number of messages was being lost during heavy overload conditions when messages began to take circuitous routes; and third, when the routing doctrine was improved to prevent such message losses, it became economically unfeasible to run the computer simulation long enough to determine the number of messages which could be expected to be lost. A target of less than one lost message per 100,000,000 was sought; this extreme requirement was selected because it was felt that future systems should be able to transmit digital data between computers--and computers are often intolerant of errors.

A new simulator, designed to resolve these problems, was encoded in the SCAT language.* This Memorandum describes that simulator and its appurtenances, and reports on the successful rectification of the previous effort's shortcomings. In addition, some analytical investigations of traffic-flow are described and evaluated.

* International Business Machines Corporation, SHARE SOS Reference Manual--SHARE Operating System for the IBM 709, IBM Applied Programming Publication, New York, 1960-61.

SUMMARY

Results of investigations into the behavior of distributed communications networks under various loading conditions are reported. A mathematical model and a deterministic equation for predicting the distribution of message path-lengths are derived and evaluated. A SCAT-encoded* simulation program that corrects deficiencies of earlier simulations is described.

An "input-choking" doctrine, together with a short, purposeful delay of messages passing through each station (when necessary), proved to be a powerful device in preventing loss of messages within networks operating at high loading ratios. The decrease in delay and in message-flow rate caused by the doctrine was negligible.

For the networks studied, a policy of dropping messages that have traversed paths greater than twice the longest possible path between the extremities of the network, resulted in a message dropout rate of less than one in 100,000,000 when the networks were operating at normal, and even higher than normal, loadings. At low loadings there were even fewer messages dropped.

*Ibid.

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SYMBOLS

<u>SYMBOL</u>		<u>INTRODUCED ON PAGE</u>
S_i	The i^{th} station in a network, in an $N \times M$ array of stations.	1
$H_i(j,k)$	The expected number of stations to be traversed by a message originating at S_i , transmitted via link j of S_i , before being delivered to S_k .	3
$L_{j,i}$	Link j of station S_i ; links are numbered from 0 through 7, and are displayed (abstractly) clockwise around the station with $L_{0,i}$ at high noon.	3
S_o	A message's station of origin.	4
S_d	A message's addressee.	4
h	The hand-over number associated with a message; i.e., the number of stations traversed by a message in its wanderings.	4
HMAX	The maximum allowable number of traverses a message may take before being dropped from circulation.	4
$B(i,k)$	By definition, the length of the best path between S_i and S_k .	5
N_x	The number of best paths of length x .	5
x'	The maximum best-path length.	5
B	The best-path distribution:	5

$$B(x) = N_x / \sum_{j=1}^{x'} N_j.$$

<u>SYMBOL</u>		<u>INTRODUCED ON PAGE</u>
λ	The number of links in a network.	8
α	Message-loading factor for a simulation.	8
π	The probability that a link is impaired (unavailable or busy).	9
ML	Message-unit length.	9

FIGURES

Figure		
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I. INTRODUCTION

The distributed communications network examined in this Memorandum is a collection of communication stations interconnected in a more or less regularized manner and employing a common switching doctrine. The connectivity is such that the number of bi-directional lines (or uni-directional links) between a station and its neighbors is constant over the interior of the network; lengths and bandwidth capabilities of the lines may vary over the network. Switching, or routing, is performed by choosing a "best" initial link toward a desired addressee, rather than by attempting to choose a "best" overall path.

Certain stylized connectivities may be introduced to define the redundancy level of such networks--the concept is illustrated in Fig. 1.* It is useful and instructive to deal with such connectivities, and we shall do so. However, many of the observations reported in this Memorandum are dependent on a constancy of connectivity and routing doctrine rather than on any specific regularity of connectivity.

The choice of a "best" initial link is effected by adjoining to each station, S_i , a matrix, H_i , which assigns

*Figure 1 is taken from ODC-I, where the concept of redundancy levels is fully described; ODC is an abbreviation of the series title, On Distributed Communications. The number following refers to the particular volume within the series.

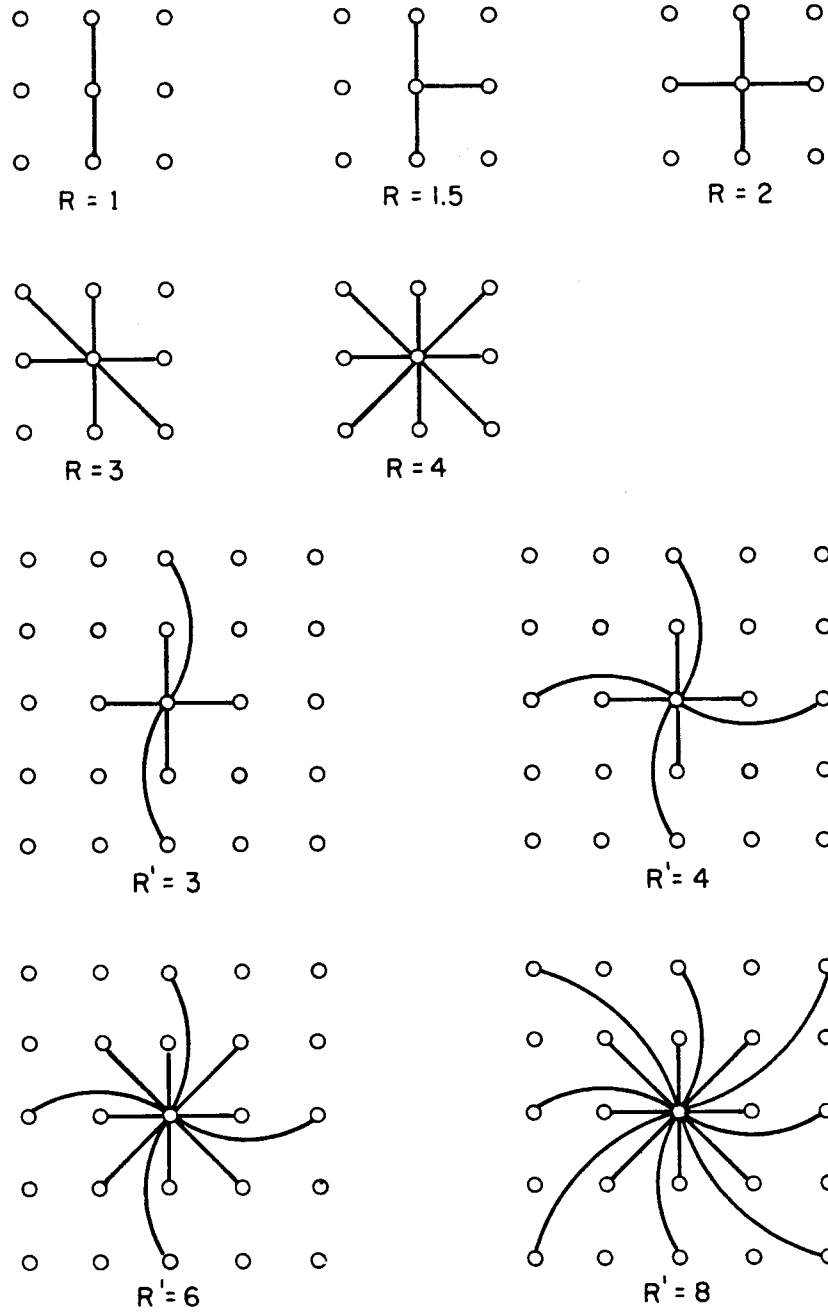


Fig. 1 Definition of Redundancy Level

to each link, $L_{j,i}$, of station S_i a measure, $H_i(j,k)$, of link j 's merit as an initial link along which to route messages destined for station S_k --for all stations, S_k , in the network. The figure of merit used in this Memorandum is the expected number of stations to be traversed before delivery of the message. That is, a value, $K = H_i(j,k)$, states that station S_i expects a message destined for station S_k that is sent out along link $L_{j,i}$ to traverse K stations before delivery. By adjoining to every message an integer, h --the hand-over number, which is incremented by unity each time the message is retransmitted by a station--the values of H_i may be continuously updated as a function of the hand-over numbers associated with incoming messages and of the current values of H_i . Messages are routed by applying a decision procedure to those H_i values pertinent to the messages' destinations.

The transient behavior of the H_i , and of message-flow as the H_i change under suitable updating algorithms, is discussed in ODC-II. The present Memorandum is concerned with "steady-state" flow, wherein the H_i have assumed their "best" values and there remain fixed.

II. THE ROUTING DOCTRINE

Let S_o , S_d , and h refer to a message's originating station, addressee, and hand-over number, respectively. Messages arriving at station S_i via link $L_{j,i}$ cause the values of H_i to be updated by the following updating doctrine (U1):

$$(U1) \quad H_i(j,o) = \text{minimum}(h, H_i(j,o)).$$

Let station S_i have links $L_{1,i}, L_{2,i}, \dots, L_{n_i,i}$ and let HMAX, the maximum allowable hand-over number, be a fixed parameter. Station S_i retransmits messages by applying the following routing doctrine (R1):

- (R1)
- a) if $h \geq \text{HMAX}$, the message is dropped;
 - b) if all links are in use or otherwise unavailable, the message is stacked until some link is freed;
 - c) if links are available, h is incremented by unity and the message retransmitted over that link $L_{j',i}$ for which

$$H_i(j',d) = \text{minimum}(H_i(j,d)), \text{ for all available links, } L_{j,i}.$$

In practice, the H_i are originally set equal to HMAX; as messages move through the network, these values eventually assume their absolute minima. These minimum values may be calculated a priori by a best-path* algorithm.

*See Sec. VI.

Assume this has been done, resulting in values \bar{H}_i . Define, for all stations S_i and S_k , $k \neq i$:

$$B(i,k) = \text{minimum}(\bar{H}_i(j,k)), \text{ for all links } L_{j,i} \text{ of station } S_i; \quad (2.1)$$

$$N_x = \text{the number of } B(i,k) \text{ equal to } x \text{ for all } i,k; \quad (2.2)$$

$$x' = \text{maximum}(B(i,k)), \text{ for all } i,k. \quad (2.3)$$

That is, $B(i,k)$ is the length of the best path from S_i to S_k , N_x is the number of such best paths of length x , and x' is the length of the longest best path. The distribution, B , defined by

$$B(x) = N_x / \sum_{y=1}^{x'} N_y \quad (2.4)$$

is called the best-path distribution for the network.

It will be shown later that B is of great utility in predicting traffic-flow in steady-state networks. If messages are generated by choosing origins and destinations randomly from a uniform distribution of stations, and if each message is completely routed through the network before the next message is originated, then the resulting distribution of message path-lengths is clearly B . That

B is equally valid for low-load conditions is shown by Fig. 3,* to be discussed later.

The HMAX test in routing doctrine R1 is essential. Traffic will, in practice, be cut into message blocks of fixed length and be sent piecemeal (but serially). If the differential time delay is excessive, then the order of arrival of these message blocks could fall out of sequence. This differential time delay is equivalent to the difference in hand-over numbers of the arriving message blocks. Therefore, the lower the tolerable maximum hand-over number, the higher the overall data rate between two network users will be. We also chose to drop message blocks whose hand-over number exceeds this maximum allowable, to insure flushing out message blocks addressed to nonexistent stations. Since fixing the length of message blocks implies fixing station processing time (for routing), network time may be equated to station traverses by using an average link-length for the network, and temporal decisions can then be made on the basis of some maximum hand-over number and the hand-over number associated with messages. Thus, the problem of choosing an HMAX small enough to maintain "integrity" of communications, yet large enough to guard against excessive message dropout is a central one. The choice

*See p. 14.

of such a value for a given network depends on a knowledge of the distribution of message path-lengths of delivered messages under varying traffic densities and using an unbounded HMAX. Three methods for determining such distributions--Monte Carlo simulation, mathematical modeling, and approximate calculation--are described and compared in the next section.

III. COMPUTATIONAL TECHNIQUES

MONTE CARLO SIMULATION

The simulator is described in detail in the program listing of Appendix B; briefly, it operates in the following manner:

- 1) a network is defined in terms of its size, configuration, and connectivity;
- 2) link-lengths are assigned--as transmission times--by being drawn from a uniform distribution bounded by the parameters TPMAX and TPMIN;
- 3) the hand-over number tables, H_i , are preset to either \bar{H}_i or to values bounded by HMAX and HPRIME;
- 4) a fixed number, $(\alpha) \cdot (\lambda)$, of messages is introduced into the network, origins and destinations of messages being drawn from a uniform distribution of station numbers*-- λ is the total number of links (twice the number of lines) in the network and α is a loading factor;
- 5) the routing and transmission of messages through the network is directly simulated by applying the routing doctrine R1;
- 6) delivery or dropping of a message results in the insertion of a new "random" message into the network, thus maintaining loading;

*This defines "uniform" loading.

7) new messages may be "choked" by applying R1 first to enrout messages, then to new messages (which are, in effect, kept on a special stack).

The parameter α is a measure of the activity of the network. We may assume that under uniform loading, traffic density can be equated with a probability, π , that a link is unavailable or busy, and may then relate π to activity by

$$\pi \approx \left[\frac{ML}{\frac{TPMAX-TPMIN}{2} + 2(ML)} \right] \cdot (\alpha) \quad (3.1)$$

where ML is the processing time per message-unit (i.e., message-unit length). $(TPMAX-TPMIN)/2$ is the average line-length (or average transmission time), and a message-unit requires ML time units to be inserted into or withdrawn from a link. In (3.1), α is treated as the number of message-units that may be on a link (either in transit or being transmitted into the link), and the first factor is interpreted as the probability that a message is being transmitted into a link (thus making the link unavailable).

Under uniform loading in a fully loaded network, at each time period all links would be jammed and all stations would be simultaneously accepting and transmitting a message. The average number of messages in a fully loaded network would, therefore, be equal to

$$(\lambda) \left(\frac{TPMAX-TPMIN}{2(ML)} + 2 \right) = N .$$

The loading ratio of a network is then equal to $N/\alpha\lambda$, which is equal to π ; thus, π may be equated to the loading ratio, or traffic density, of a network.

MODEL A (M_a)

A model of network behavior is abstracted by considering the four possibilities that exist for the disposition of messages at any station (disregarding the possibility of dropout). A message being routed at a station has associated with it that station's prediction of the best-path length to the message's addressee. The message may either

- 1) remain at the station, or
- 2) be retransmitted to a neighbor whose best-path prediction is
 - a) one (1) less than the previous station's prediction,
 - b) the same as the previous,
 - c) one (1) greater than the previous.

That is, links may be characterized as being "best," "next best," or "worst" for any message. The model defined in Fig. 2 assumes:

- 1) a uniform distribution of links with respect to "best," "next best," and "worst" overall addressees;

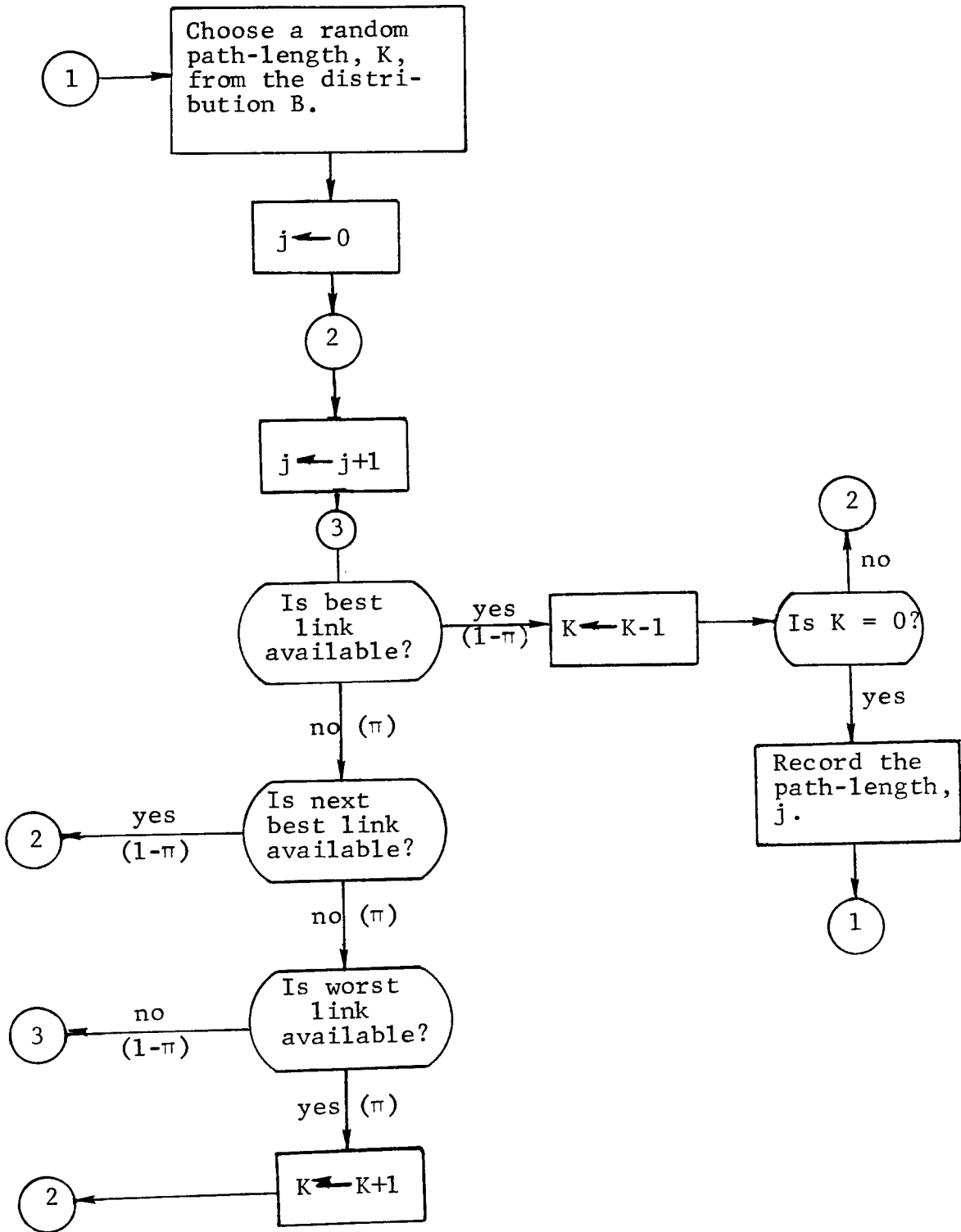


Fig. 2 Flowchart of Model A

- 2) a uniform probability, π , of link nonavailability;
- 3) uniform loading;
- 4) stations connected to neighbors only.

The first assumption is the weakest of the four; however, within the interior of "reasonably" connected networks satisfying the last assumption, it is "reasonably" valid.

MODEL B (M_b)

Model B is obtained from the previous one by aggregating "next best" and "worst" links. In M_b , therefore, a message at a station whose best-path prediction for the message is, say, n , will have a probability $p(n,x)$ of being delivered in x traverses, given by

$$p(n,x) = g(n,x) (\pi^{x-n}) (1 - \pi)^n, \quad (3.2)$$

where π is the probability of link nonavailability, and $g(n,x)$ is the number of combinations of "best" paths (chosen with probability $1 - \pi$) and "worst" paths (chosen with probability π) ending in a "best" path.

That is:

$$\begin{aligned} g(n,x) &= \begin{bmatrix} x-1 \\ n-1 \end{bmatrix} \\ &= g(n,x-1) + g(n-1,x-1). \end{aligned} \quad (3.3)$$

The distribution, $M(\pi)$, generated by this model is given by:

$$M(x;\pi) = \sum_{i=1}^{x'} B(i) \cdot P(i,x) \quad . \quad (3.4)$$

Letting $\pi = 1/a$, and using (3.2), we obtain:

$$M(x;1/a) = \sum_{i=1}^{x'} B(i) \cdot (a-1)^i \cdot (a)^{-x} \cdot g(i,x) \quad . \quad (3.5)$$

Since $\sum_{x=1}^{\infty} g(i,x) \cdot (a)^{-x} = (a-1)^{-i}$, we have:

$$\sum_{x=1}^{\infty} M(x;1/a) = \sum_{i=1}^{x'} B(i) = 1 \quad ; \quad (3.6)$$

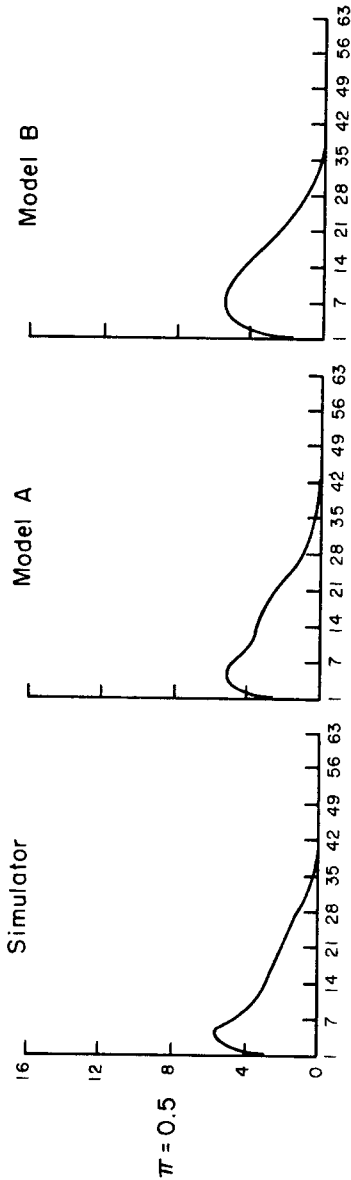
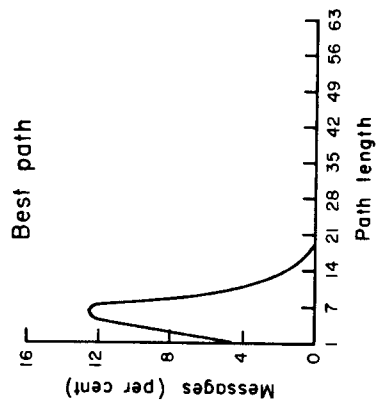
and for the "tail" of the distribution from $x'+k$ on:

$$\begin{aligned} \sum_{x=x'+k}^{\infty} M(x;1/a) &= \sum_{i=1}^{x'} B(i) \cdot (a-1)^i \left(\sum_{x=x'+k}^{\infty} g(i,x) \cdot (a)^{-x} \right) \\ &< \sum_{i=1}^{x'} B(i) \cdot (a-1)^{i-x'-k} \quad . \end{aligned} \quad (3.7)$$

COMPARISON OF RESULTS

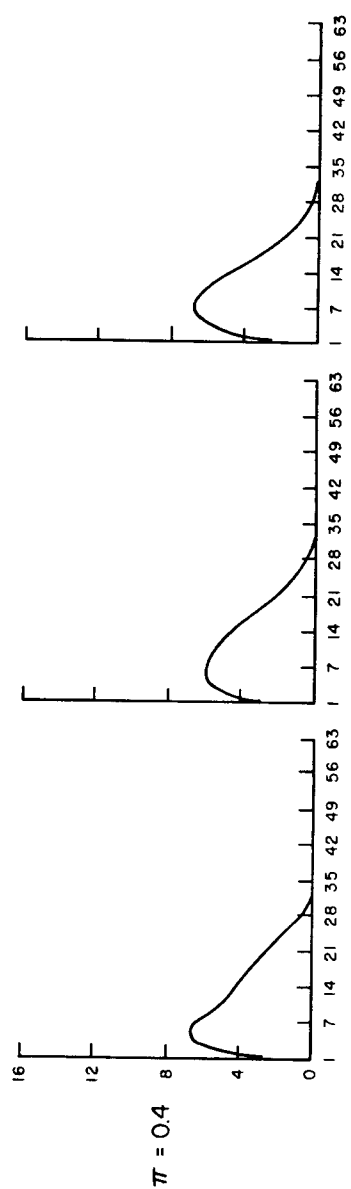
The three techniques were applied to networks of size 10x10 and 14x7, using redundancy levels of 2, 3, and 4. Figure 3 exhibits the results of the 14x7, redundancy-three case--all other runs produced the same behavior; Fig. 4 tabulates statistics for the various runs.

π = Probability that each link is busy or impaired.



Model A

Model B



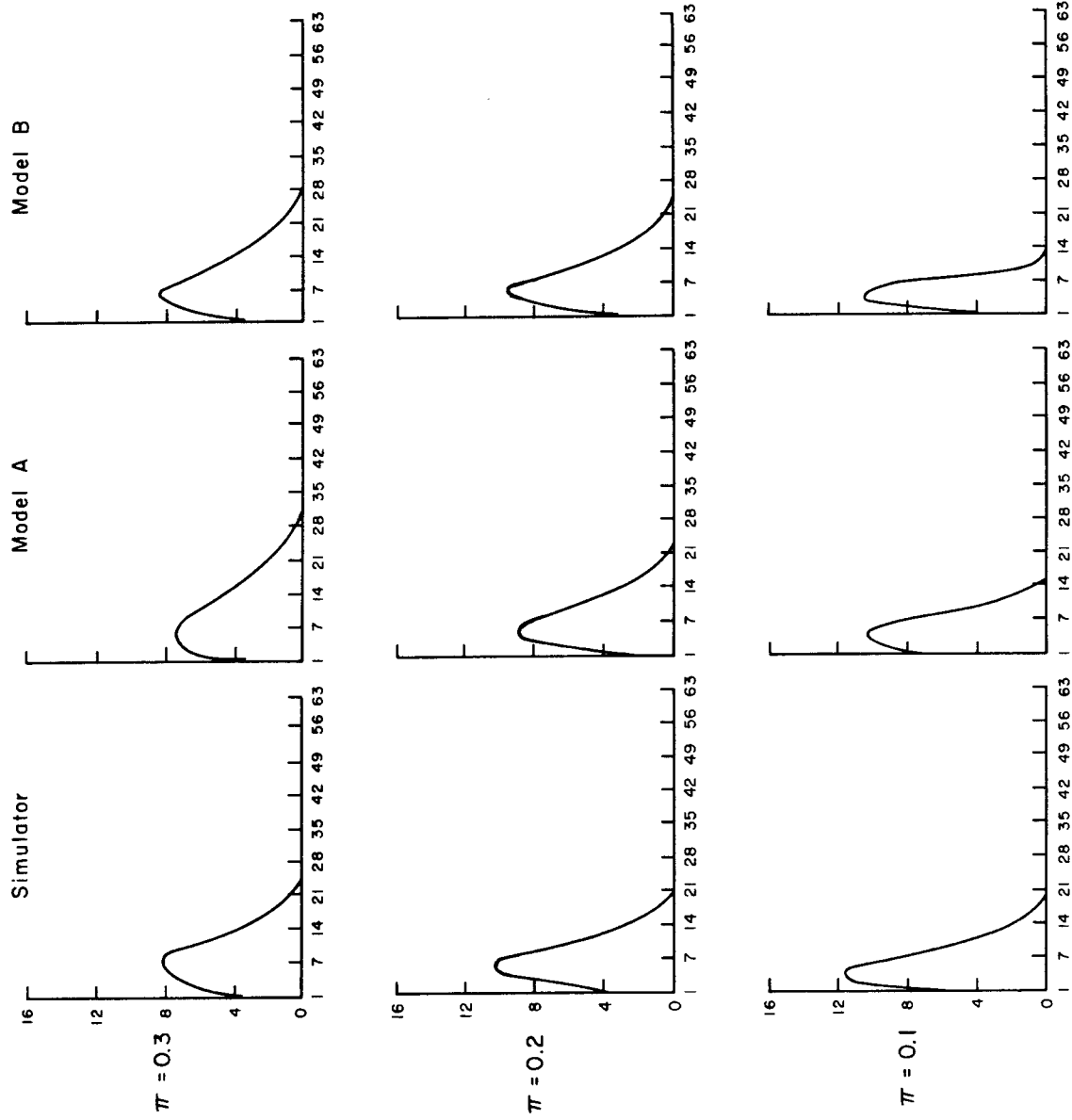


Fig. 3 Distribution of Path-Lengths of Delivered Messages:
14 x 7 Network; R = 3

Relation (3.1) was used to equate simulator loading with the probability of link impairment, π , used by the models. Since this is an approximate relation, no exact comparisons between simulation and modeling can be made. Nevertheless, it is clear that both models produce distributions which are commensurate with those produced by the simulator, Model A tending to reproduce the shape of the simulator distributions more faithfully than Model B. Moreover, note that both models tend to produce distribution "tails" which are generally pessimistic. The same comparisons held for the other runs. We therefore conjecture that both models may be used to predict approximate behavior of networks of the type examined, and that (3.7) may be used to obtain approximate estimates of HMAX's for desired dropout rates.

Figure 4 indicates that the transition from a redundancy level of three to a level of four results in diminishing returns. For redundancy levels greater than two, it appears that an HMAX equal to twice the longest best-path is sufficient to reduce dropouts to the "noise" level under normal, and even higher than normal, loadings ($\leq 30\%$).

IV. INPUT-CHOKING AND STACK LENGTHS

Input-choking* in networks has two important consequences. First, since a link can be usurped by a message for no greater period of time than it takes to insert a message into a link (ML), it is clear that stack length can never exceed the number of links associated with a station; that is, stations require only one "word" of stack storage per link. Moreover, choking tends to smooth out potential activity peaks, particularly in the vicinity of "hog" stations and around the center of the network. Removal of input-choking simply means that new messages are treated as enroute messages. Under such conditions, stack lengths cannot be contained. Since stack storage must be finite, a new message-dropping criterion must be added to the routing doctrine, R1:

- (R1) d) if no links are available and the stack is full, the message is dropped.

Thus, the choice of a fixed stack-storage length, STACK, large enough to reduce stack dropouts to the noise level, becomes as important as a choice of HMAX. Since stack length is a highly local phenomenon--strongly dependent on traffic distributions--there seemed to be no simple way of determining STACK without direct simulation. Accordingly, the simulator was applied--with no input-choking--to a 14x7 net of redundancy-three, with maximum stack

*See p. 9.

lengths of 6, 9, and 12. The distributions obtained were almost identical to the "choked" distributions. They are not reproduced here; however, Fig. 5 compares the pertinent simulations, and leads us to the following conjectures:

1) dropouts will occur under no-choke conditions; to insure a low dropout rate, STACK will have to be unconscionably large--hence, unfeasible;

2) under uniform loading, no-choking produces a small, uniform increase in loading;

3) hence, there seems to be no justification for adopting a no-choking doctrine.

These conjectures are fortified by a consideration of the distribution of total time in stacks (including input-waiting time) for delivered messages. These distributions were identical for the choke and no-choke cases. In fact, the distributions shown in Fig. 6 were typical of all networks simulated; they show that the probability of excessive delays-in-stack is exceedingly small, even under high loading.

Probability of Link Impairment, π	HO Average	HO Variance	Mode Value of HO No.	Choke?	Stack Drops	No. in Stack
.1	6.5	13.2	4	yes	0	12
	6.5	13.5	4	no		
	6.5	13.5	4	no		
.3	8.0	20.6	5	yes	0	12
	8.1	21.7	5	no		
	8.1	21.7	5	no		
.5	13.7	104	3-7	yes	5	12
	14.3	117	3-7	no		
	14.2	114	3-7	no		
.7	17.9	166	3-7	yes	27	12
	17.4	167	3-7	no		
	17.9	176	3-7	no		

Fig. 5 No-Choking Statistics

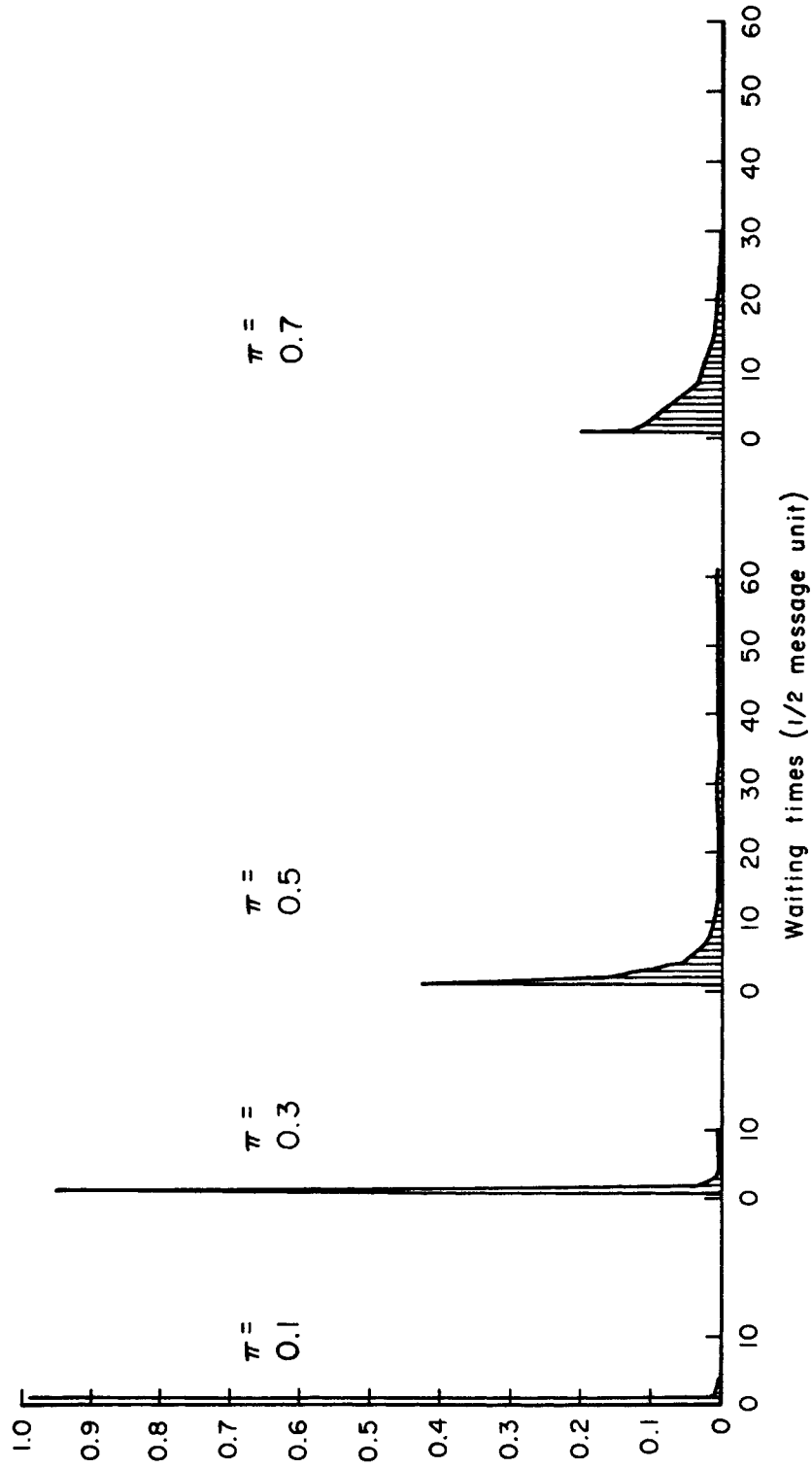


Fig. 6 Distribution of Stack Waiting Times of Delivered Messages
(14 x 7 Network; Redundancy 3; Choking)

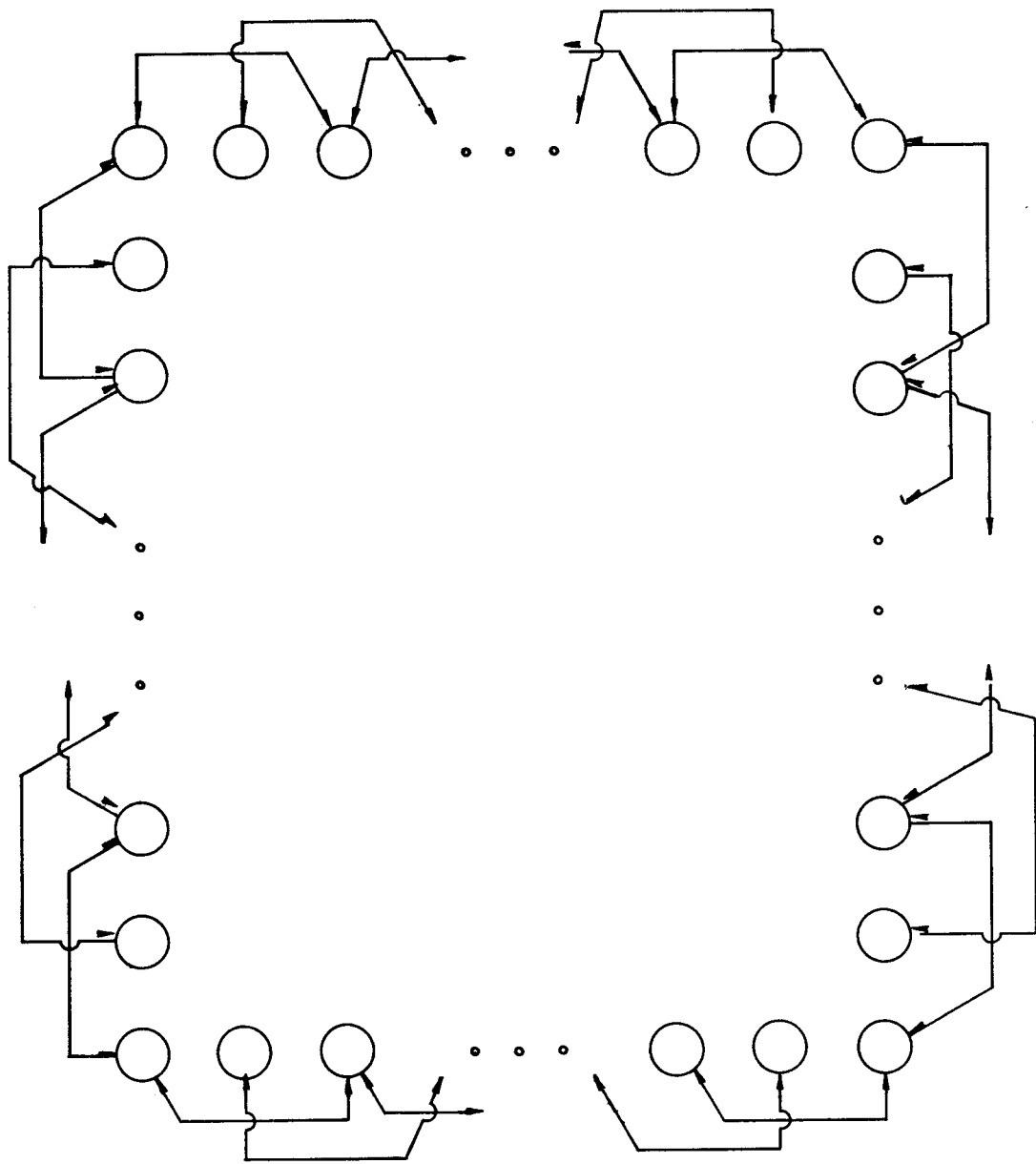


Fig. 7 Edge Binding

V. EDGE BINDING

In any network, much of the routing power of peripheral stations is wasted simply because peripheral links are unused. Thus, messages tend to reflect off the boundary into the interior or to move parallel to the periphery. Providing alternate paths by edge binding, as illustrated in Fig. 7, should tend to shorten average path-lengths measurably. Figure 8 exhibits flow patterns produced by simulations of a 14x7 network of redundancy-three with and without edge binding. Each diagram is a spatial representation of the network, the entry in position (i,j) of the representation indicating the number of messages routed through the corresponding station since the start of the simulation. The effect of edge binding in reducing interior clustering is clear. Figure 10 exhibits the distributions for the edge-bound net, while Fig. 9 gives statistics for the various runs. Although edge binding reduces clustering and results in an increase in flow-rate, it seems clear that the resultant distributions are less desirable than those obtained without edge binding. Apparently edge binding tends to overly penalize paths through the center of the network, and seems to provide higher flow rates at the expense of slightly higher drop-out rates.

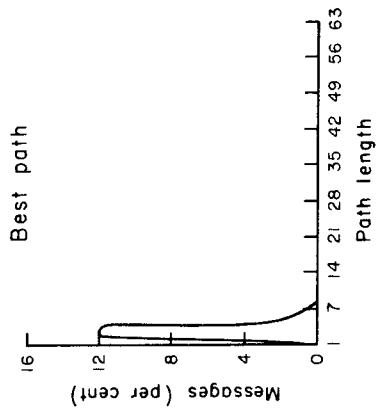
	UNBOUND							BOUND						
	219	408	514	601	566	525	413	872	919	1893	2452	4131	3715	4641
	506	941	1209	1419	1510	1689	1375	1664	1326	1344	2097	3277	4859	4742
	759	1397	1866	2403	3110	3582	2788	2942	2238	1787	2408	3763	5125	5790
	1045	2043	2914	4066	4773	5128	3567	4135	3194	2714	3216	4284	5365	5862
	1476	3114	4399	5263	5579	5687	3860	4753	3971	3451	3831	4615	5477	5836
	2256	4481	5328	5702	5858	5876	3890	5273	4472	4031	4260	4861	5454	5827
	3036	5326	5705	5875	5920	5868	3819	5477	4784	4234	4471	4916	5352	5777
$\pi=0.5$	3513	5656	5824	5855	5864	5695	3575	5548	4807	4322	4429	4749	5296	5757
	3699	5678	5798	5770	5653	5286	3071	5478	4826	4231	4165	4489	5066	5671
	3689	5620	5634	5476	5089	4191	2109	5498	4592	3719	3633	3886	4732	5487
	3568	5282	5145	4613	3710	2657	1251	5312	4208	3156	2855	3098	4000	5111
	3002	4114	3641	2923	2161	1561	799	5168	3732	2464	2045	2051	2972	4077
	1732	2072	1796	1457	1223	1028	551	4022	3279	1859	1428	1203	1768	2434
	493	579	645	627	567	477	249	3747	2240	2701	1442	1444	880	1270
	287	512	594	638	597	490	330	567	616	1179	1216	2109	1763	3138
	658	1122	1300	1434	1404	1296	760	889	776	842	1124	1328	2402	3332
	937	1621	1939	2206	2198	1929	1165	1512	1125	1169	1523	1795	2710	4578
	1176	2096	2554	2933	2895	2538	1500	2448	1592	1674	1974	2113	3195	4796
	1444	2542	3070	3466	3374	2995	1791	3118	2210	2013	2239	2410	3472	5124
	1664	2893	3491	3833	3701	3353	2037	3934	2500	2203	2425	2663	3683	5153
$\pi=0.3$	1842	3224	3782	4054	3932	3553	2120	4170	2850	2325	2454	2670	3669	5109
	2028	3484	3889	4072	3950	3471	1969	4521	2961	2387	2466	2599	3537	4883
	2129	3442	3717	3893	3718	3099	1718	4521	3084	2327	2413	2453	3168	4631
	2002	3153	3448	3518	3292	2659	1462	4558	2857	2155	2137	2143	2742	3952
	1678	2689	2942	2999	2731	2160	1160	4212	2531	1849	1933	1732	2015	3129
	1284	2018	2277	2262	2070	1665	904	3876	2040	1494	1424	1213	1345	2064
	827	1350	1477	1445	1357	1111	631	2725	1873	1180	1042	836	830	1149
	424	542	636	646	593	489	273	2655	1492	1910	1171	1149	634	656
	116	205	244	239	233	183	115	251	279	555	484	891	531	1255
	250	466	542	568	538	480	274	355	334	344	432	431	579	966
	371	665	808	819	790	687	404	624	472	511	625	599	618	1606
	430	800	932	1007	1006	874	509	924	548	738	751	714	690	1637
	516	967	1081	1189	1151	943	591	1122	673	752	860	797	760	1994
	556	988	1169	1298	1216	998	627	1452	675	804	867	832	785	1872
	591	1058	1273	1307	1281	1052	643	1515	701	791	918	861	748	1998
$\pi=0.1$	598	1066	1274	1344	1280	1072	599	1755	666	817	885	830	794	1748
	570	1030	1179	1276	1183	977	577	1571	717	819	886	827	675	1618
	533	959	1089	1150	1086	938	511	1707	691	819	869	781	685	1297
	510	826	949	994	981	757	441	1362	629	723	772	702	602	1028
	415	637	780	810	776	604	346	1423	583	602	606	570	500	646
	287	483	520	520	522	387	261	782	527	407	417	355	323	419
	142	177	216	214	224	180	113	1089	487	805	462	520	269	251

**Fig. 8 Traffic Flow Distribution
(Number of Messages Routed Through Node in 7 x 14 Network)**

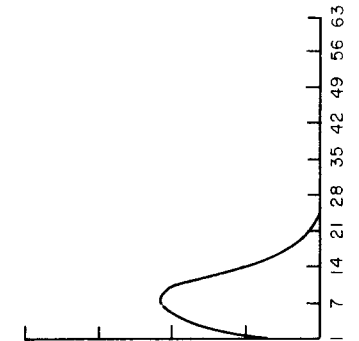
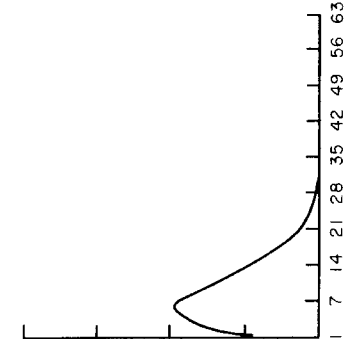
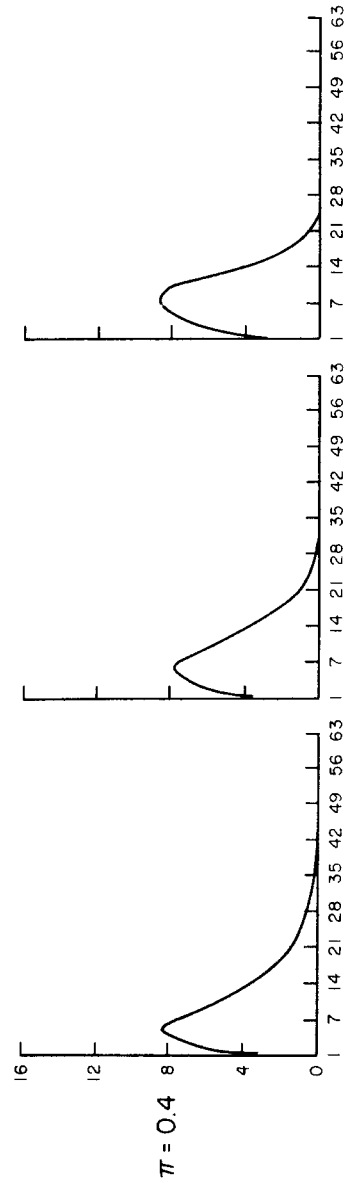
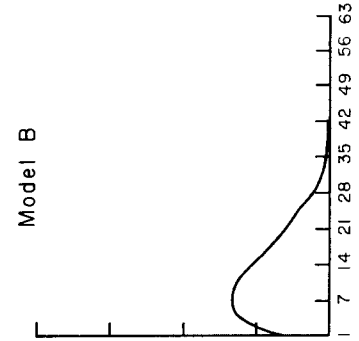
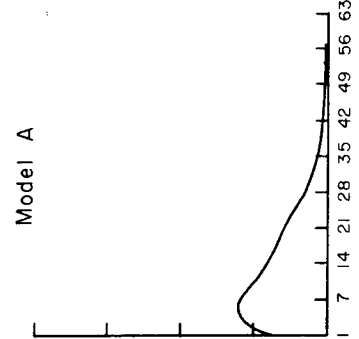
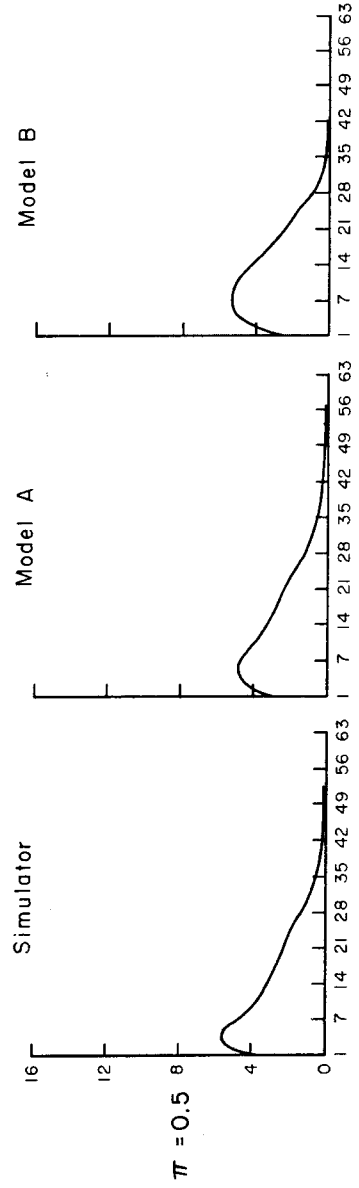
Net Description	Best-Path		π	Average Path-Length			HMAX Required to Reduce Dropouts to One in 10^8 .		
	Average	Longest		S_1	M_a	M_b	S_1	M_a	M_b
14 x 7 R = 4	4.32	10	.5	9.98	10.10	8.61	>63	61	52
			.4	7.77	7.98	7.18	59	44	40
			.3	6.29	6.58	6.16	38	36	32
			.2	5.35	5.58	5.39	23	23	25
			.1	4.78	4.87	4.79	15	16	20
14 x 7 R = 3	4.86	10	.5	12.10	11.31	9.70	>63	>63	55
			.4	9.85	9.02	8.09	>63	59	42
			.3	7.84	7.43	6.93	50	31	34
			.2	6.34	6.27	6.07	31	24	27
			.1	5.44	5.48	5.39	17	18	21
14 x 7 R = 2	5.46	11	.5	14.11	12.66	10.89	>63	>63	57
			.4	11.20	10.09	9.07	62	60	44
			.3	8.56	8.36	7.78	41	36	35
			.2	7.00	7.06	6.81	32	27	28
			.1	6.01	6.11	6.05	20	17	22
10 x 10 R = 3	4.83	10	.5	10.48	11.27	9.65	>63	>63	54
			.4	8.10	8.97	8.04	44	59	42
			.3	6.74	7.37	6.89	31	31	33
			.2	5.84	6.23	6.03	21	23	26
			.1	5.26	5.44	5.36	17	17	20

S_1 - Simulation M_a - Model A M_b - Model B

Fig. 9 Statistics for Networks: Edge Binding



π = Probability that each link is busy or impaired.



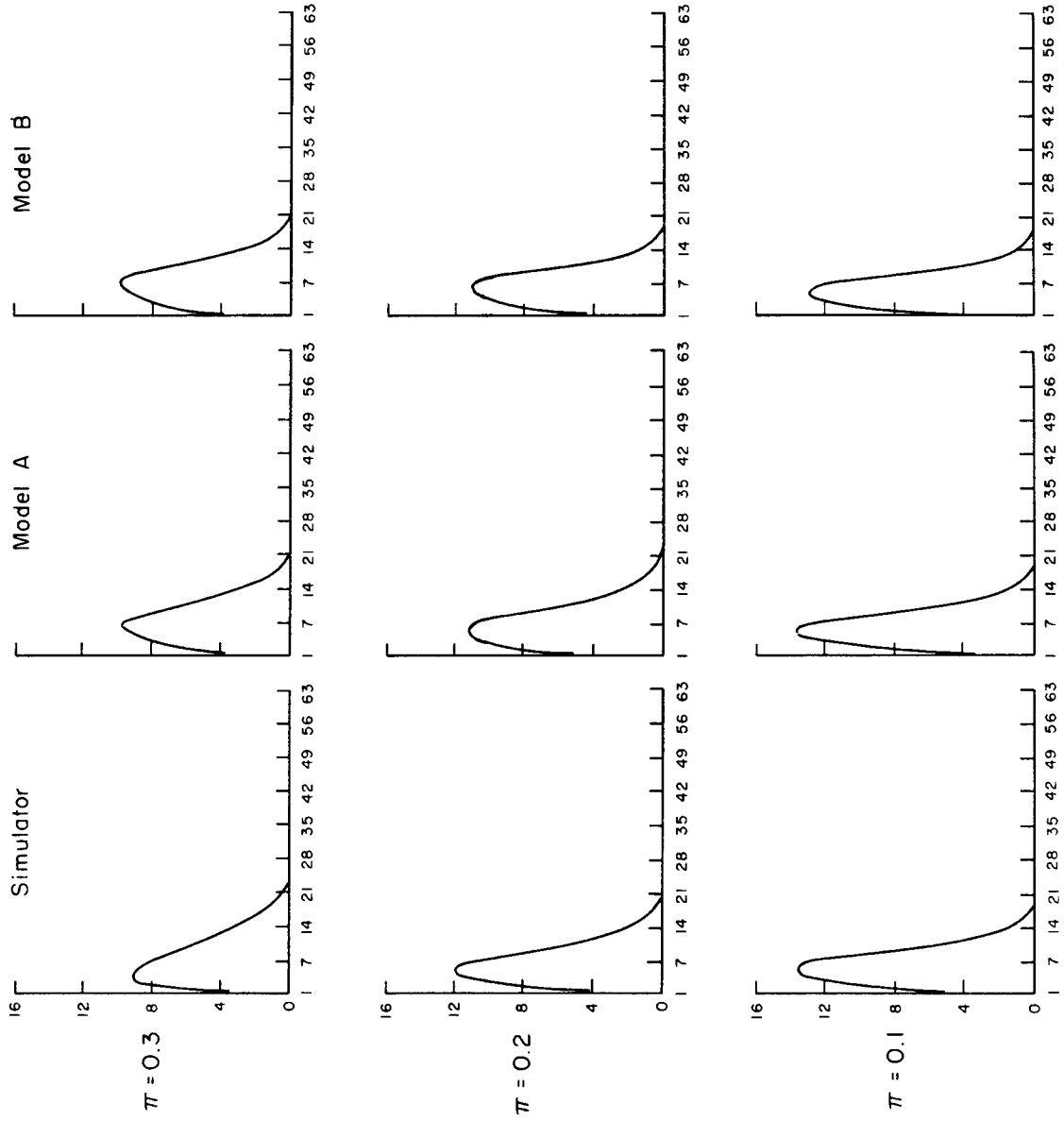


Fig. 10 Distribution of Path-Lengths of Delivered Messages:
14 x 7 Network; R = 3; Edge Binding

VI. VARIATIONS IN TRAFFIC DENSITY

Techniques for smoothing out fluctuations in traffic density are discussed in ODC-II. A technique suggested and evaluated in that Memorandum uses a modification of the HO-table updating algorithm which allows the values of $H_i(j,k)$ to adapt to changing traffic conditions. The adaptation algorithm is described in the footnote to Fig. 12 of the present Memorandum (p. 37), which considers the effects of lengthy bursts of activity from a single station or between a pair of stations. Such effects can be studied by appending to the simulator provisions for treating specified stations as message sources or sinks, with specified weights associated with each station so chosen. A station specified as a source (sink) of weight $k < 1$ will be chosen by the simulator as a message originator (addressee) with probability k . Since correct receipt of each message block is acknowledged by the adjacent receiving station, it is clear that any prolonged increase in the rate of message generation by a single station will result in a commensurate increase in network loading. In such cases we would expect effects similar to those produced by an increase in message loading. In cases where many stations suddenly increase their rate of message generation to a particular station, a similar increase in network loading should occur. This

is prevented from occurring by mechanisms described in ODC-VIII.

Several hog cases were simulated on a 7x7 network of redundancy three. The procedure used was the following: first, simulate the network at a 20 per cent loading ratio without sources or sinks; then increase the loading ratio to 40 per cent, specify a commensurate source/sink configuration, and continue the simulation, comparing statistics at each step. The results are summarized:

- 1) Regions of heavy flow activity centered about sources and sinks;
- 2) Source/sink configurations caused a 50 per cent decrease in the message-flow rate--this was probably caused by input-choking;
- 3) Distributions of stack-waiting times--with sources and sinks--were identical to those obtained for 40 per cent loadings without sources or sinks;
- 4) Single stations acting as both source and sink had little effect on the distribution of path-lengths of delivered messages, the distributions being only slightly less desirable than those obtained for 20 per cent loading with no sources or sinks;
- 5) Two stations, each acting as both source and sink, had effects on the path-length distribution which depended on the relative positions of the stations within

the network--the more remote the stations, the worse the distributions.

We conjecture that single stations acting as a hog source will reduce message-flow rate but will result in few, if any, dropped messages.

The case of a single source was simulated, such as would occur when a "fraudulent" station attempted to overload the network. The results were as anticipated--the increased loading produced distributions expected from the new loadings. Single sources had essentially the same effects as single source/sinks, with the exception that stack waiting times remained relatively unchanged.

VII. THE BEST-PATH ALGORITHM

Although the algorithm used to set the hand-over number tables, H_i , to their best values was designed solely to minimize computer running time, it might also conceivably find use other than in the Distributed Adaptive Message Block Network in allowing broadcast of best-path information through a network. The algorithm requires a specific, recognizable message-type, an info-message, and a variation of the standard routine doctrine to process such messages. The algorithm is as follows:

- 1) a station, S_o , that wishes to broadcast best-path information to the rest of the network originates an info-message and transmits the message over all its links, $L_{j,o}$;
- 2) a station, S_i , receiving an info-message, characterized by (S_o, h) , via link $L_{j,i}$ compares h with $H_i(j,o)$:
 - a) if $h \geq H_i(j,o)$ the message is dropped;
 - b) otherwise, $H_i(j,o)$ is set to h , h is incremented by unity, and the message is retransmitted over all of S_i 's links.

The best-path algorithm used is a parallel application of 1) and 2) above, with all stations acting as originators once and only once. In practice, info-messages may be characterized either by the absence of an addressee or by a "universal" addressee.

Appendix A

PROGRAM DESCRIPTION

A listing of a collection of SCAT-encoded computer routines designed to operate under the aegis of a user-composed supervisory routine is contained in Appendix B. The collection includes routines for defining networks in terms of pertinent parameters, for assigning and re-assigning parameter values, for performing Monte Carlo simulations on networks, for applying Model A and Model B to networks, and for displaying results of simulations and model-runs. The routines are operative on the IBM 7090, require the RAND versions of SOS* for that machine, and are well-described by the listing of Appendix B.

Supervisory routines to perform network calculations must be encoded in SCAT and should use the macro-directives described on p. 4 of the listing. Figure 11 contains an example of a supervisory routine, suitably annotated. The general procedure is to first assign parameter values, then to simulate or model, and finally to display or interrogate results and, perhaps, iterate the procedure. Network parameters are described in Fig. 12; the "normal" values there listed remain in effect until changed by the

* Bryan, G. E., Ed., The RAND-SHARE Operating System Manual for the IBM 7090 Computer, The RAND Corporation, RM-3327-PR, September 1962.

```

*   JOB      8109,TEST,JWS618,70,35000,35,C
*   ASSIGN   B2=SYSBR1
*   LOAD     GOIF,NOSQZ,NOLIST,SYSBR1
*   CHANGE   NETPGM      BEGINNING OF SUPERVISORY ROUTINE.
*   ROW      14          14X7 NETWORK,
*   COLUMN   7
*   WEAVE    3           REDUNDANCY THREE,
*   CHOKE    WITH INPUT CHOKING, AND
*   NOBIND   NO EDGE-BINDING.
*   FLOW     MESSAGE FLOWS DESIRED.
*
*   FIRST, PRINT BEST-PATH DISTRIBUTION.
*   BESTHO   SYSTEM      EXIT IF NOGO.
*   THEN APPLY SIMULATOR, MODEL A, AND MODEL B -- VARYING LOADING FROM 10
*   PERCENT TO 50 PERCENT IN JUMPS OF 10, USING UNBOUNDED HMAX(NORMAL CASE).
*
.A   AXT      5,1
*   IMPAIR   .L,1        SET LOADING, THEN
*   SIMUL    SYSTEM,3000 SIMULATE FOR 3000 CYCLES. THEN DISPLAY
*   PRINT    S1DIST      PATH-LENGTH DISTRIBUTION,
*   PRINT    S1WAIT      STACK-WAITING-TIME DISTRIBUTION, AND
*   PRINT    S1FLOW      MESSAGE-FLOW PATTERNS.
*   MODELA   SYSTEM      NEXT, APPLY MODELS. ( DISTRIBUTIONS
*   MODELB   SYSTEM      ARE DISPLAYED).
*   TIX      .A,1,1
*   IMPAIR   .L2
*   SIMUL    SYSTEM,3000 NEXT RUN FOR 3000 CYCLES AT
*   IMPAIR   .L4          20 PERCENT LOADING.
*   SOURCE   1,.5        NEXT, DOUBLE LOADING AND DEFINE STATION
*   CONT     1000        1 AS A SOURCE OF WEIGHT 1/2.
*   PRINT    S1DIST
*   PRINT    S1WAIT
*   PRINT    S1FLOW
*   SINK     1,.5        CONTINUE, THEN DISPLAY.
*
*   CONT     1000
*   PRINT    S1DIST
*   PRINT    S1WAIT
*   PRINT    S1FLOW
*   TRA      SYSTEM      FINI
*   DEC      .1
*   .L2     DEC      .2
*   .L4     DEC      .3
*   .L4     DEC      .4
*   .L4     DEC      .5
*   .L      EQU      *
```

Fig. 11 Sample Network Program

supervisory program. Much of Figs. 11 and 12, and the macro-directive listing in Appendix B, is self-explanatory.

The parameter WEAVE, which specifies network connectivity, is defined by Fig. 13. The parameter GRAIN defines the ratio of message-unit length (time required to insert a message into a link) to the time required by a station to route a message. Since message-routing time is equivalent to simulation cycle time, GRAIN defines the "coarseness" of the simulation. Note that simulations may be halted and then continued; during these "pauses" the user may display results of the simulation and may change the values of certain parameters. Parameters which may be changed during the course of a simulation are indicated in Fig. 12 with an asterisk. Sources and sinks may be defined, deleted, or have their weights changed during a simulation. Deletion of a source or sink is accomplished by assigning a zero weight. One further restriction exists: the impairment factor, IMPAIR, may not be reset to a value greater than that which held at the initiation of the simulation; this, however, is not a real restriction, since the simulator can be initiated for a zero-time run.

Execution times for application of the two models are negligible. Simulation times are a function of the size and connectivity of the net being simulated and of the traffic density. A reasonable approximation to

Parameter Name	Function	Restrictions	Normal Value
*CHOKE	= 0 implies no input-choking; ≠ 0 implies input-choking.	none	≠ 0
ALPHA	The loading factor, described in Sec. III.	computed as function of IMPAIR	0
HMAX	The maximum hand-over number; HMAX = 63 implies messages are never dropped.	HMAX an integer; $0 < HMAX \leq 63$	63
HPRIME	Used to preset the hand-over table, H; HPRIME = 0 implies H is preset to "best" values, otherwise H is preset with values drawn from uniform distribution between HPRIME and HMAX.	h' an integer $0 \leq h' \leq HMAX$	0
ROW	rows	integers	--
COLUMN	Number of columns		
GRAIN	The "grain" of the simulation; the ratio of message-length (time to place message on or accept message from a link) to the time required by a station to process the message.	integer > 0	2
TPMAX	Link-lengths are assigned values drawn from uniform distribution bounded by (TPMAX+GRAIN, TPMIN+GRAIN).	integral number of station processing times	6
TPMIN			0

Fig. 12 Network Parameters

Parameter Name	Function	Restrictions	Normal Value
STACK	Maximum stack-storage, described in Sec. IV.	integer	8
WEAVE	The connectivity of the network; see Fig. 13.		--
*IMPAIR	The probability of link impairment; used by M_a and M_b , and the simulator.	decimal fraction	0
RANDOM	An initial pseudo-random number.	large, odd integer	976525005
FLOW	$\neq 0$ implies message-flow statistics are desired; $= 0$ implies none are desired.	none	$\neq 0$
BIND	$\neq 0$ implies edge binding (Fig. 7); $= 0$ implies no edge binding.	none	0
*ADAPT	$= 0$ implies H_i table entries, h_i , are updated by $h_i = \text{minimum}(h_i, h)$; $\neq 0$ implies H_i table entries are updated by adaptation algorithm. ^a	none	0
*LEARN	Learning factor for adaptation.	$0 < \text{LEARN} \leq 1$	1
*FORGET	Forgetting factor for adaptation.	$0 < \text{FORGET} \leq 1$	0

^aLet $D = (h - h_i)$

$>$ $H_0 = H_0 + D \cdot \text{FORGET}$
 If $D = 0$ then $H_0 = H_0$
 $<$ $H_0 = H_0 + D \cdot \text{LEARN}$

Fig. 12 (Continued)

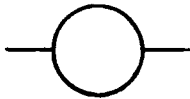
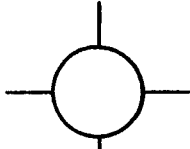
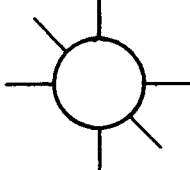
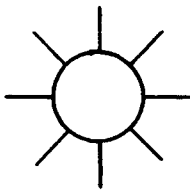
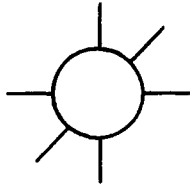
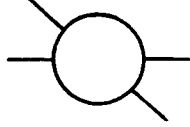
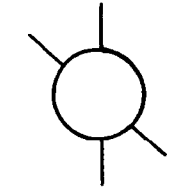
<u>WEAVE</u>	<u>Redundancy Level</u>	
1	1	
2	2	
3	3	
4	4	
5	3	
6	2	
7	2	

Fig. 13 WEAVE Parameter

execution times for an $n \times m$ net of redundancy r is given by:

$$\text{time per simulation cycle} = \frac{2 \cdot r \cdot n \cdot m \cdot \alpha}{3} \text{ ms}$$

where α is the loading factor described in Sec. III.

Computer storage required is a function of net size and connectivity, of traffic density, and of other desiderata. An $n \times m$ net of redundancy- r requires approximately

$$(4 \cdot n \cdot m)(1+r+\alpha) + \frac{(n \cdot m)}{2} \cdot [n \cdot m \cdot r - (r-1)(n+m)]$$

words of computer storage, 24,000 words being available. A 10x10 network of redundancy-four, or a 11x11 network of redundancy-three can be accommodated.

Appendix B

PROGRAM LISTING

In this appendix is presented the program listing, together with a Table of Contents to the routines listed. The page numbers referred to in this table are the numbers internal to the listing, not the pages in the Memorandum.

TABLE OF CONTENTS

* INDICATES ROUTINES OF DIRECT INTEREST TO USERS

ROUTINE	DESCRIPTION	PAGE
MACRO DESCRIPTIONS		6
NETDMP	ERROR DUMPS AND DISPLAYS	7
NETPGM	SUPERVISORY PROGRAM	7
PARSET	USED TO SET PARAMETER VALUES	8
.ENTRY	ENTRY ROUTINE FOR SUPERVISORY MACROS	10
PRESET	INITIALIZES NETWORKS. USED BY SIMULATOR AND BY MODELS. USES INL, INS, IA, IL, BH.	11
INL	SETS CONNECTIVITY MAP	16
INS	COMPUTES STORAGE MAP	17
IA	SETS LINK TABLE	21
IL	GENERATES LINK LENGTHS	23
IH	GENERATES INITIAL HO TABLE	26
BH	GENERATES BEST HO TABLE	29
MCVER	USED BY SIMULATOR TO MOVE MSGS	31
HCVER	USED BY MOVER AND BH TO UPDATE HO TABLE	34
UPDATE	* USED BY HOVER AS HO-TABLE UPDATING ALGORITHM	36
DELIVR	USED VARIOUSLY TO STACK MESSAGES AT STATIONS	37
RCUTER	* USED BY SIMULATOR TO ROUTE MSGS. USES FA, FB	39
FA	DETERMINES A STATIONS AVAILABLE LINKS	42
FB	ORDERS AVAILABLE LINKS W.R.T. HO-TABLE VALUES	43

DROP	PLACES DROPPED MSGS ON DROP LIST	44
RECCRD	PLACES DELIVERED MSGS ON ARRIVAL LIST	45
CHAIN	USED VARIOUSLY TO PLACE MSGS ON LISTS	47
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PARAM	NETWORK PARAMETERS	77
TABLES	LISTS AND TABLES USED BY VARIOUS PROGRAMS	77

LIST OF MACRO-DIRECTIVES

ALL MACROS PRESERVE IR'S AND SENSE

'I' MEANS EITHER OF 1) INTEGER
2) LOCATION
3) LOCATION TAG

MACROS FOR SETTING PARAMS

BINC
NCBINC

CHOK
NCCHOK

FLOW
NCFLW

ADAPT
NCDAPT NO ADAPTATION TO BE USED

HMAX I
WEAVE I
RCW I
COLUMN I
TPMAX I
TPMIN I
STACK I
HPRIME I
GRAIN I
IMPAIR LOC,TAG

MACROS TO DEFINE 'HOG' STATIONS. STATION
N IS TO BE ASSIGNED AS A MESSAGE
SOURCE (OR SINK) WITH PROBABILITY K

SOURCE N,K K A DECIMAL FRACTION
SINK N,K

MACROS TO DEFINE SPECIAL LINKS TO BE ADDED
OR DELETED.

LINK N1,K1,N2,K2 LINK FROM STATION N1 LINK
NR. K1, TO STATION N2, LINK
NR K2, IS TO BE ADDED.

CLT N1,K1,N2,K2

MACRC DIRECTIVES

'E' MEANS 'ERROR EXIT'
 'L' MEANS 'LOCATION OF SPECIAL-LINK CONTROL-WORD'
 WHERE (L) CONTAINS --
 LOCATION OF FIRST SPECIAL LINKS, NR OF SPECIAL LINKS
 AND L=0 OR (L)=0 IMPLIES NO LINKS
 I REFERS TO EITHER LOCATION OR INTEGER

SIMUL	E, I, L	START S1 SIMULATION, RUN FOR I OR (I) CYCLES
CCNT		CONTINUE SIMUL FOR I OR (I) CYCLES
MDELA	E, L	APPLY MODEL A. DISPLAY DIST.
MDELB	E, L	APPLY MODEL B. DISPLAY DIST.
BESTHO	E, L	GENERATE HO-TABLE BY BEST-PATH ALGORITHM. DISPLAY 'BEST' DIST., HBAR
SETHMX	X, L	USED IMMEDIATELY AFTER APPLICATION OF MODEL A OR B TO SET HMAX TO A VALUE THAT, ON THE BASIS OF THE APPLICABLE MODEL'S RESULTS(DISTRIBUTION) , WILL INSURE A FRACTIONAL DROP-OUT NOT EXCEEDING (L) EXITS TO X IF NO SUCH VALUE EXISTS. MAY ALSO BE USED AFTER OR DURING SIMULATION RUN TO ADJUST HMAX.
PRINT	S1DIST	DIST. OF PATH-LENGTHS OF MSGS DLVRD BY S1
PRINT	S1FLOW	TRAFFIC FLOW GENERATED BY S1
PRINT	S1WAIT	DIST. OF STACK WAITING TIMES

BEGINNING OF NETWORK PROGRAM.
NETPGM TSX FORMAT,4 TO DEFINE XFORMATS
SUPERVISORY PROGRAM FOLLOWS
+1 TRA SYSTEM PROTECTION

DUMP CLA* 1,4
+1 TZE 2,4
+2 ALS 18
+3 ADD* 1,4
+4 STO *+5
+5 STL SYSCB1
+6 TXL SYSCB2,,21
CORE LIMIT,LIMIT,C
+7 STL SYSCB1
+11 TRA 2,4

PANEL
ERROR STL SYSCB1
+2 TSX DUMP,4
+3 PZE TABLES
+4 TSX DUMP,4
+5 PZE LINTBL
+6 TSX DUMP,4
+7 PZE BUSY
+8 TSX DUMP,4
+9 PZE NCDTBL
+10 TSX DUMP,4
+11 PZE MSGTBL
+12 TSX DUMP,4
+13 PZE HODTBL
+14 TSX DUMP,4
+15 PZE HOTBLE
+16 TSX DUMP,4
+17 PZE STDTBL
+18 TSX DUMP,4
+19 PZE SCTBLE
+20 TSX DUMP,4
+21 PZE SKTBLE
+22 TSX DUMP,4
+23 PZE POOL
+24 TSX DUMP,4
+25 PZE PARAM
+26 TSX SIDIST,4
+27 TSX SIFLOW,4
+28 TSX SIWAIT,4
+29 TRA SYSTEM

PARSET USED TO SET ALL PARAMS

```

CALL  SEQU  STL  .EXIT
        TXL  PARSET,,PARAM NAME
        TXL  LOC,TAG,0 1 OR 2

```

SAVES IR'S AND SENSE

```

PARSET  SXA      PRSX,4
+1  CLA*      .EXIT
+2  ARS       18
+3  STA      PRS1          PARAM NAME
+4  CLA      .EXIT
+5  ADD      K.A1
+6  STA      PRS2
+7  ADD      K.A1
+8  STA      .EXIT
PRS2    CLA      **0          FETCH ARGUMENTS
+1  STA      PRS4          LOC OF NEW VALLE(CR VALUE ITSELF)
+2  STT      PRS4
+3  LDQ*     PRS4
+4  PCX      ,4          C,1 OR 2
+5  TXH     PRS3,4,0      SETTING FLOATING PCINT QUANTITIES
+6  CAL     SYSORG      ELSE TEST IF
+7  ANA     K.M6          VALLE OR LOC OF VALUE
+8  CAS     PRS4
+9  LDQ     PRS4          VALLE
+10 LDQ     PRS4          VALUE
+11 XCA
+12 SSP
+13 CAS     K.M2
+14 TRA     **3
+15 TRA     **2
+16 TRA     PRS1-1
+17 XCA
+18 PXA
+19 LLS     8
+20 SUB     KINT2
+21 TPL     **2
+22 PXA
+23 STA     **2
+24 PXA
+25 LLS     **0
+26 XCA
PRS1    STQ     **0          LOC OF VALLE
+1  TRA     PRSX
PRS3    XCA
+1  SSP
+2  CAS     K.M2
+3  TRA     **4
+4  NOP
+5  ORA     KINT1
+6  FAD     KINT1

```


	+7	STC*	PRS1
PRSX		AXT	**0,4
	+1	TRA*	.EXIT
PRS4		PZE	

.ENTRY USED FOR ALL MACRO DIRECTIVES
 CALL SEQU STL .EXIT
 TXL .ENTRY,,LENGTH OF CALL SECU + 2
 TSX PROCESSOR,4
 CALLING SEQU FOR PROCESSOR

MAX CALL SEQU LENGTH = 8

.EXIT	PZE	**0	
.ENTRY	SWT	5	EXCESS TIME TEST
+1	TRA	**2	CONTINUE
+2	TRA	SYSTEM	FINI
+3	CLA	.EXIT	
+4	ALS	18	
+5	STD	.EN1	
+6	SXA	.ENX,4	
+7	CLA*	.EXIT	
+8	PDX	,4	
+9	TXI	**+1,4,.EN2	
+10	SXA	.EN3,4	
+11	PDX	,4	
.EN1	TXI	**+1,4,**C	
+1	SXA	.EN4,4	
+2	SXA	.EXIT,4	
+3	PDX	,4	
+4	TXI	**+1,4,-1	
.EN4	CLA	**0,4	
.EN3	STO	**0,4	
+1	TIX	.EN4,4,1	
+2	CLA	.EN5	
+3	AXT	,4	
.EN2	STO*	.EN3	
+1	TSX	ERROR,4	
+2	TSX	ERROR,4	
+3	TSX	ERROR,4	
+4	TSX	ERROR,4	
+5	TSX	ERROR,4	
+6	TSX	ERROR,4	
+7	TSX	ERROR,4	
+8	TSX	ERROR,4	
+9	TSX	ERROR,4	
+10	TSX	ERROR,4	
.EN5	TRA	.ENX	
.ENX	AXT	**0,4	
+1	TRA*	.EXIT	
.ERROR	AXT	1,4	
+1	CLA*	.EN3	
+2	ARS	18	
+3	STA	**2	
+4	LXA	.ENX,4	
+5	TRA	**0	

PRESET VALIDATES PARAMS, DISPLAYS PARAMS
INITIALIZES NET, AND ALLOCATES STORAGE.

CALL SEQU TSX PRESET,4
PZE LOC OF FIRST SPEC. LINK,,NR OF LINKS
RETURN IF STORAGE EXCEEDED CR ERRCRS
NORMAL RETURN

PRES	STL	DSPY	ENTRY FROM SIMLLATCR
+1	TRA	PRO	
PRESET	STZ	DSPY	ENTRY FROM MODELS AND BEST-PATH DSPLY
	BEGIN	3,7,1	
PRO	TXL	++7,0,0	SUBROUTINE LINKAGE
+11	CLA	1,4	
+12	STO	PR5	
		SET RANDOM AND SET ACTIVITY LEVELS	
+13	CLA	TPMAX	
+14	SUB	TPMIN	
+15	SSP		
+16	ORA	KINT1	
+17	FAD	KINT1	
+18	SUB	KINT3	(TPMAX-TPMIN)/2 = D
+19	STO	K.T1	
+20	CLA	GRAIN	
+21	ORA	KINT1	
+22	FAD	KINT1	ML = MESSAGE LENGTH
+23	STO	K.T2	
+24	FAD	K.T2	
+25	FAD	K.T1	
+26	FDP	K.T2	(D+2ML)/ML = ACTIVITY FACTOR
+27	STQ	FACTOR	
+28	FMP	IMPARE	ACTIVITY LEVEL
+29	STO	ALPHA	
+30	STO	MAXACT	
+31	CLA	RANDOM	
+32	LBT		
+33	CLA	K.R	
+34	SLW	RANDOM	
+35	SLW	RND	
+36	NZT	DSPY	
+37	TRA	PR7	
		DISPLAY PARAMETERS	
+38	TSX	PARAMS,4	
+39	TRA	PR7	

PARAMS PRINTS PARAMETER VALUES.

	BEGIN	1,7,1	
PARAMS	TXL	++7,0,0	SUBROUTINE LINKAGE

```

      XEJECT
+11 STL SYSOED
+14 AXT 4,2
+15 AXT RANDOM-ALPHA,1
      XPRINT D,2
PR1.1 STL SYSOED
      +3 PZE PM,1,1
      +4 PZE RANDOM,1,1
      +5 TXI *+1,1,-1
      +6 TIX PR1.1,2,1
      XPRINT E,2
PR3   STL SYSOED
      +3 PZE PM,1,1
      +4 PZE RANDOM,1,1
      +5 TIX PR3,1,1
      XPRINT F,2
      +6 STL SYSOED
      +9 PZE PM,,1
+10 PZE RANDOM,,1
+11 AXT 0-4,2
PR6   AXT ,1
      +1 ZET CHOKE-1,2
      +2 AXT 1,1
      XPRINT G,2
      +3 STL SYSOED
      +6 PZE PMC-1,2,1
      +7 PZE NO,1,1
      +8 TXI *+1,2,1
      +9 TXH PR6,2,0
      RETURN PARAMS
+10 TRA PARAMS+1

```

```

      FLUSH OUTPUT BUFFERS
PR7   STL PTTGL
      +1 TSX SYSDSK,4
      +2 PZE SYSMOT
      +3 TSX SYSDSK,4
      +4 PZE SYSMIT
      +5 STZ SYSIBC
      +6 TSX SYSBFD,4
      +7 PZE SYSSBF,,512
      +8 STR
      +9 STZ BFTST
      TEST AND SET CONNECTIVITY
+10 LXA WEAVE,4
+11 TXH **2,4,0
+12 TRA PRX
+13 TXL **2,4,7
+14 TRA PRX
+15 LAC WEAVE,4
+16 CLA CTABLE,4
+17 STO LINKS
      TEST LINK-LENGTH LIMITS

```

+18	CLA	TPMAX	
+19	TNZ	*+2	
+20	TRA	PRX	
+21	CAS	TPMIN	
+22	TRA	*+3	
+23	TRA	*+2	
+24	TRA	PRX	
+25	ADD	GRAIN	
+26	ADD	K.A1	
+27	CAS	TPHIGH	
+28	TRA	PRX	
+29	NOP		
		TEST HANDOVER-NUMBER LIMITS	
+30	CLA	MAXHO	
+31	ALS	3	
+32	TNZ	*+2	
+33	TRA	PRX	
+34	CAS	H0H11	
+35	CLA	H0H11	
+36	NOP		
+37	STO	HOMAX	
+38	CLA	INITHO	
+39	ALS	3	
+40	CAS	HOMAX	
+41	TRA	PRX	
+42	NOP		
+43	STO	HOMEL	
+44	AXT	LIMIT-MSG2-1,1	
+45	STZ	LIMIT,1	CLEAR STORE
+46	TIX	*-1,1,1	
+47	LXD	CNTTBL,1	CLEAR COUNTS
+48	STZ	CNTTBL,1	
+49	TIX	*-1,1,1	
+50	AXT	TPU,1	CLEAR TRAFFIC-LIST HEADS
+51	STZ	TP,1	
+52	TIX	*-1,1,1	
+53	SXA	TP,1	
+54	TSX	INL,4	SET LINK TABLE
+55	TSX	INS,4	GENERATE STORAGE MAP
+56	TRA	PRX	DRAT
+57	NZT	NODES	TEST VALIDITY OF PARAMS
+58	TRA	PRX	
+59	NZT	LINES	
+60	TRA	PRX	
+61	NZT	MSGS	
+62	TRA	PRX	
+63	TSX	IA,4	SET LINE TABLE
+64	TSX	IL,4	SET LINK-LENGTHS AND LINKS
PR5	PZE	**0,,**0	DEFINES SPECIAL LINKS
+1	TRA	PRX	
+2	TSX	IH,4	
+3		HOMAX	
+4		HOMEL	
+5	TRA	PRX	

```

                NOW CALCULATE EXACT LCADING FROM TRUE NR OF LIMKS
+6 LDQ      NODES
+7 MPY      LINKS
+8 STQ      K.T4
+9 LDQ      HONCNT
+10 PXA
+11 DVP     NODES
+12 XCA
+13 STO     K.T3
+14 XCA
+15 MPY     MSGS
+16 DVP     K.T4
+17 CAL     MSGTBL
+18 ARS     19
+19 ANA     K.M6
+20 TLQ     **2
+21 XCA
+22 STQ     MSGS
+23 STQ     LOAD
+24 CLA     HONCNT
+25 SUB     K.T3
+26 STO     HONCNT
                EXACT LOADING
                ALLOT EXTRA BUFFERS IF STORAGE AVAILABLE
+27 LAC     IH7,1
+28 TXI     *+1,1,32767
+29 TXL     PR4,1,255
+30 STL     BFTST
+31 SXD     **4,1
+32 LXA     IH7,1
+33 SXA     **2,1
+34 TSX     SYSBFD,4
+35 PZE     **0,,**0
+36 STR
RETURN PRO
PR4 TRA PRO+1
   DUMP TABLES
PRX TSX DUMP,4
   DUMP PARAM
+2 TSX DUMP,4
RETURN PRO,1
+4 AXT 1,4
PMC BCI 1,CHOKE
+1 BCI 1,FLOW
+2 BCI 1,BIND
+3 BCI 1,ADAPT
+4 BCI 1,ALPHA
+5 BCI 1,IMPAIR
+6 BCI 1,LEARN
+7 BCI 1,FORGET
+8 BCI 1,HMAX
+9 BCI 1,WEAVE
+10 BCI 1,ROW
+11 BCI 1,COLUMN
+12 BCI 1,TPMAX

```

```
+13 BCI 1,TPMIN  
+14 BCI 1,STACK  
+15 BCI 1,HPRIME  
+16 BCI 1,GRAIN  
PM BCI 1,RANDOM
```

INL PRESETS LINK TABLE AS FCT OF COLUMN DIM.

CALL SEQU TSX INL,4
RETURN

INL	TRA	*+1	
	BEGIN	1,7	
INLO	TXL	*+5,0,0	SUBROUTINE LINKAGE
+8	CLA	CX	
+9	STA	INL1	
+10	PAC	,1	
+11	AXT	4,4	
+12	AXT	1,2	
INL3	SXD	INL2,2	
+1	PXA	,1	
+2	STA	LINK+4,4	
+3	PAC	,2	
+4	PXA	,2	
+5	STA	LINK+8,4	
INL2	TXI	*+1,1,**C	
INL1	AXT	**0,2	
+1	TIX	INL3,4,1	
	RETURN	INLO	
+2	TRA	INLO+1	

INS COMPUTES STORAGE MAP AS FUNCTION OF

CX = COLUMN DIM
 RX = KCW **
 LINKS = LINKS PER NODE
 ACTIVE = PERCENT MSG ACTIVITY

CALL SEQ TSX INS,4
 RETURN IF STORE EXCEEDED
 NORMAL RETURN

```

INS   TRA   **1
      BEGIN 2,7
INSO  TXL   **5,0,0          SUBROUTINE LINKAGE
      +8 LDQ  CX
      +9 RQL  3
      +10 STQ D1             8*COLUMN DIM
      +11 MPY  RX
      +12 STQ LINES         8*NR OF NODES = LINE TABLE LENGTH
      +13 LRS  3
      +14 STQ NODES         NR OF NODES
      CALCULATE MSG LOADING USING APPARANT NR OF LINKS
      +15 MPY  LINKS
      +16 XCA
      +17 ORA  KINT1
      +18 FAD  KINT1
      +19 XCA
      +20 CLA  ALPHA
      +21 CAS  K.M2
      +22 TRA  **3
      +23 NOP
      +24 STZ  ALPHA
      +25 FMP  ALPHA
      +26 SSP
      +27 XCA
      +28 PXA
      +29 LLS  8
      +30 SUB  KINT2
      +31 TPL  **2
      +32 PXA
      +33 STA  **2
      +34 PXA
      +35 LLS  **0
      +36 ZET  ACTTGL
      +37 PXA          NO MSG STG NEEDED FOR MODELS
      +38 ADD  K.A1
      +39 STO  MSGS     DESIRED NR OF MSGS
      NEXT ALLOCATE STORAGE FOR ALL TABLES
      +40 ADD  RX       ALLOW EXTRA STORAGE
      +41 ADD  RX       FOR POSSIBLE BINDING
      +42 ADD  CX       OF COARSELY CONNECTED
      +43 ADD  CX       NETS, OF RED 2 FOR EXAMPLE.
      +44 ALS  1       LENGTH OF MESSAGE TABLE
      +45 PAX  ,1
  
```

+46	SXD	MSGTBL,1	
+47	ADD	K.A1	
+48	ADD	BASE	TOP OF LINE TABLE
+49	STO	LINTBL	
+50	LXA	LINES,1	
+51	SXD	LINTBL,1	
+52	ADD	LINES	
+53	STA	LINE1	
+54	ADD	K.A7	BASE OF LINE TABLE
+55	STA	LINEA	TOP OF NODE TABLE
+56	STA	LINEB	
+57	STA	LINEC	
+58	ADD	K.A1	
+59	STO	NCDTBL	
+60	CLA	NCDES	
+61	CAS	K.M5	
+62	TRA	INSX	OVERSIZE NET
+63	TRA	INSX	
+64	STO	SKH	
+65	STO	SCH	
+66	ADD	NODTBL	
+67	STA	N1A	
+68	STA	N1B	
+69	STA	N1C	
+70	ADD	NCDES	
+71	ADD	K.A1	
+72	STA	N2A	
+73	STA	N2B	
+74	STA	N2C	
+75	ADD	NCDES	
+76	ADD	K.A1	
+77	STA	N3A	
+78	STA	N3B	
+79	STA	N3C	
+80	NZT	SNAP	
+81	TRA	*+6	
+82	ADD	NCDES	
+83	ADD	K.A1	
+84	STA	N4A	
+85	STA	N4B	
+86	STA	N4C	
+87	SUB	NODTBL	
+88	PAX	,1	
+89	ADD	NCDTBL	
+90	SXD	NODTBL,1	
+91	ADD	K.A1	
+92	STO	HODTBL	
+93	ADD	MAXHO	
+94	STA	HODA	ALLOCATE HO-DISTRIBUTION TABLE
+95	STA	HODB	
+96	STA	HODC	
+97	ADD	K.A1	
+98	STO	STDTBL	
+99	ADD	MAXHO	

```
+100 STA STDA
+101 STA STDB
+102 STA STDC
+103 LXA MAXHD,2
+104 SXD HODTBL,2
+105 SXD STDTBL,2
+106 ADD K.A1
+107 STO SCTBLE
+108 ADD SSLIM
+109 STA SC1A
+110 STA SC1B
+111 STA SC1C
+112 ADD K.A1
+113 ADD SSLIM
+114 STA SC2A
+115 STA SC2B
+116 STA SC2C
+117 ADD K.A1
+118 STO SKTBLE
+119 ADD SSLIM
+120 STA SK1A
+121 STA SK1B
+122 STA SK1C
+123 ADD K.A1
+124 ADD SSLIM
+125 STA SK2A
+126 STA SK2B
+127 STA SK2C
+128 STZ SOURCE
+129 STZ SINK
+130 STZ SKCUM
+131 STZ SCCUM
+132 LXA SSLIM,1
+133 SXD **1,1
+134 TXI **1,1,**0
+135 SXD SCTBLE,1
+136 SXD SKTBLE,1
+137 ADD K.A1
+138 STO BUSY
+139 ADD LINES
+140 ADD K.A7
+141 STA BUSYA
+142 STA BUSYB
+143 STA BUSYC
+144 LXA LINES,1
+145 SXD BUSY,1
+146 ADD K.A1
+147 STO HOTBLE
+148 CAS UPPER
+149 TRA INSX
+150 TRA INSX
+151 CAS LOWER
+152 TRA **3
+153 TRA INSX
```

+154	TRA	INSX
+155	LXA	NCDES,1
+156	PXD	,1
+157	ADD	NODES
+158	STO*	N3A
+159	SUB	K.A101
+160	TIX	=-2,1,1
+161	CLA	LINE1
+162	ADD	D1
+163	SUB	LINES
+164	STA	BOTTOM
+165	ADD	LINES
+166	SUB	K.A8
+167	STA	LEFT
	RETURN	INSO
+168	TRA	INSO+1
	RETURN	INSO,1
INSX	AXT	1,4

PRESET REFERENCE TABLE

IA INITIALIZES LINE TABLE,
IMAGINARY LINES ARE MADE BUSY.

CALL SECL TSX IA,4
RETURN
CCNN = CONNECTIVITY TYPE
CTABLE(...) DESCRIBED ELSEWHERE

IA	TRA	**+1	
	BEGIN	1,7	
IA0	TXL	**+5,0,0	SUBROUTINE LINKAGE
+8	LXA	CONN,1	
+9	LXA	LINES,2	
IA1	AXT	8,4	
+1	LDC	CTABLE,1	
+2	PXA		
IA2	LLS	0	BUSY-FLAG OFF FOR REAL LINE ON FOR IMAG LINES
+1	STC*	LINE1	
+2	RQL	1	
+3	TNX	IA3,2,1	
+4	TIX	IA2,4,1	
+5	TRA	IA1	
IA3	LXA	D1,2	NOW SET BOUNDARIES
+1	SXD	IA9,2	
+2	CAL	K.MZE	
IA4	LDQ	B.T.	
+1	AXT	8,4	
IA5	TQP	**+2	
+1	ORS*	BOTTOM	
+2	RQL	18	
+3	TQP	**+2	
+4	ORS*	LINE1	TCP
+5	RQL	19	
+6	TNX	IA6,2,1	
+7	TIX	IA5,4,1	
+8	TRA	IA4	
IA6	LXA	LINES,1	
IA7	SXA	**+1,1	
+1	AXT	**0,2	
+2	AXT	8,4	
+3	LDQ	SIDES	
IA8	TQP	**+2	
+1	CRS*	LINE1	RIGHT
+2	RQL	18	
+3	TQP	**+2	
+4	CRS*	LEFT	
+5	RQL	19	
+6	TNX	IA9,4,1	
+7	TXI	IA8,2,-1	
IA9	TIX	IA7,1,**C	
+1	LXA	LINES,1	DO SAME FOR BUSY TABLE
+2	TXI	**+1,1,7	
+3	CLA*	LINEA	

```
+4 STO* BUSYA
+5 TXI  *+1,1,-1
+6 TXH  *-3,1,7
      RETURN  IAO
+7 TRA  IAO+1
```

IL INSERTS NEW LINKAGES, ASSIGNS TP'S, DRAWN
FROM UNIFORM DIST. BTWN TP-BCUNDS, TO ALL
REAL LINES, AND POINTS IMAGE LINES AT EACH OTHER.

CALL SEQU TSX IL,4
PZE A,,B
ERROR RETLRN
RETURN

WHERE B = NR OF SPECIAL LINKS
A = ADDRESS OF FIRST LINK SPECIFIER.

LINKS FROM (I1,J1) TO (I2,J2) ARE SPECIFIED BY

LINK I1,J1,I2,J2

	BEGIN	3,7	
IL	TXL	**+5,0,0	SUBRCUTINE LINKAGE
+8	CLA	TPMAX	
+9	SUB	TPMIN	
+10	ADD	K.A1	
+11	STO	IL.T	
+12	CLA	1,4	
+13	PDX	,1	
+14	TXL	ILO,1,0	NO SPECIAL LINKS
+15	AXT	IL4,1	
+16	SXA	IA17,1	SET RETURN SWITCH.
+17	STL	K.T2	
+18	PAX	,4	
+19	SXD	**+1,1	
+20	TXI	**+1,4,**0	
+21	SXA	IL1,4	
+22	LXA	LINES,2	
+23	TXI	**+1,2,7	HIGHEST LINE NR
+24	SXD	IL2,2	
+25	SXD	IL3,2	
IL1	CLA	**0,1	GET LINK SPECIFIER
+1	PAX	,2	(I1,J1)
+2	PDX	,4	(I2,J2)
+3	TXL	ILX,2,7	ARE LINES WITHIN BCUNDS
+4	TXL	ILX,4,7	
IL2	TXH	ILX,2,**0	
IL3	TXH	ILX,4,**C	
+1	TPL	IA10	TO IA10 IF LINK
+2	CLA	K.MZE	OTHERWISE, CUT LINK
+3	STO*	LINEB	
+4	STO*	LINEC	
+5	STO*	BUSYB	
+6	STO*	BUSYC	
IL4	TIX	IL1,1,1	
ILO	ZET	BIND	TEST FOR EDGE-BINDING.
+1	TRA	IA18	TO BE BOUND.
ILOO	AXT	IA15,1	NOW, SET
+1	SXA	IA17,1	RETURN SWITCH, AND
+2	LXA	LINES,2	BEGIN ORDINARY LINKS.
+3	TXI	**+1,2,7	

	+4	STZ	K.T2	
IA14		CLA*	LINEB	
	+1	TMI	IA15	IGNORE DEAD LINES
	+2	TNZ	IA15	IGNORE PROCESSED LINES
IA10		LDQ	RND	
	+1	MPY	K.R	
	+2	STQ	RND	
	+3	MPY	IL.T	
	+4	ADD	TPMIN	
	+5	ADD	GRAIN	TP + MSG LENGTH USED AS TP
IA13		ALS	18	
	+1	STO*	LINEB	SET TP
	+2	ZET	K.T2	
	+3	TRA	IA16	SPECIAL LINE
	+4	STD	K.T1	
	+5	PXA	,2	GENERATE
	+6	LGR	3	IMAGE
	+7	PAX	,4	LINE
	+8	PXA		NR.
	+9	LGR	3	
	+10	CAQ	LINK,1,1	
	+11	SXD	**1,1	
	+12	TXI	**1,4,**0	
	+13	XCL		
	+14	PXA	,4	
	+15	LGL	3	
	+16	PAX	,4	
	+17	CLA	K.T1	
IA16		STO*	LINEC	SAME TP TO
	+1	PXA	,2	IMAGE LINE.
	+2	STO*	BUSYC	IMAGES
	+3	PXA	,4	POINT
	+4	STO*	BUSYB	TO
IA17		TRA	**0	EACH OTHER
IA15		TXI	**1,2,-1	RETURN SWITCH.
	+1	TXH	IA14,2,7	
		RETURN	IL	
	+2	TRA	IL+1	
		RETURN	IL,1	
ILX		AXT	1,4	
IA18		AXT	IA19,1	BINDING EDGES.
	+1	SXA	IA17,1	SET RETURN SWITCH
	+2	CLA	CX	COLUMN BECOMES
	+3	STO	.LS-4	TOP-EDGE-COUNT,
	+4	STO	.LS-2	BOTTOM-EDGE-COUNT,
	+5	ALS	3	
	+6	STO	.LD-3	RIGHT-SIDE DELTA, AND
	+7	SSM		SCALED
	+8	STO	.LD-1	LEFT-SIDE DELTA (-).
	+9	CLA	ROW	
	+10	STO	.LS-3	ROW BECOMES
	+11	STO	.LS-1	SIDE-COUNTS.
	+12	XCA		
	+13	MPY	.LD-3	ROW X COLUMN (SCALED)

+14	STQ	K.T1	
+15	XCA		
+16	ALS	18	
+17	ADD	K.T1	
+18	SUB	K.A16	
+19	ORA	.LK1-2	LINK(RC,4,RC-2,3)
+20	STO	.LK2-2	
+21	CLA	.LD-3	
+22	ALS	17	
+23	ADD	.LD-3	
+24	ALS	1	
+25	ADD	.LD-3	
+26	ORA	.LK1-3	LINK(C,2,3C,1).
+27	STO	.LK2-3	
+28	CLA	K.T1	
+29	ADD	K.A8	
+30	SUB	.LD-3	RC-C+1
+31	STO	K.T1	
+32	ALS	18	
+33	ADD	K.T1	
+34	SUB	.LD-3	
+35	SUB	.LD-3	
+36	ORA	.LK1-1	LINK(RC-C+1,6,RC-3C+1,5).
+37	STO	.LK2-1	
+38	AXT	4,1	INITIALIZE EDGE-COUNT
+39	TRA	IA20+2	
IA20	AXT	**0,1	
+1	TNX	IL00,1,1	FINI IF ALL EDGES BOUND.
+2	SXA	IA20,1	
+3	CLA	.LD,1	
+4	ALS	18	
+5	ADD	.LD,1	NEW INCREMENT.
+6	STO	K.T2	
+7	LDQ	.LK2,1	FIRST LINK IN NEXT EDGE.
+8	CLA	.LS,1	NEW COUNT+2
+9	PAX	,1	
+10	XCA		
+11	TXI	IA21,1,-2	
IA19	AXT	**0,1	
+1	TNX	IA20,1,1	OUT IF EDGE HAS BEEN BOUND.
+2	CLA	K.T1	
+3	ADD	K.T2	INCREMENT LINK
IA21	STO	K.T1	
+1	SXA	IA19,1	
+2	PAX	,2	
+3	PDX	,4	
+4	TRA	IA10	TO SET LINKS IN NET.

IH CHAINS REAL LINES TO HO TABLE, WHICH IS
 PRESET FROM UNIFORM DIST BETWEEN (A) AND (B).

CALL SEQU TSX IH,4
 A
 B
 RETURN IF STORE EXCEEDED
 NORMAL RETURN

	BEGIN	4,7	
IH	TXL	*+5,0,0	SUBROUTINE LINKAGE
+8	CLA	NODES	
+9	ADD	K.A3	
+10	ARS	2	NK WORDS PER HO-TABLE ENTRY
+11	STO	HOLINE	AT FOUR ITEMS PER WORD
+12	CAL*	2,4	
+13	XCA		
+14	CLA*	1,4	
+15	SSP		
+16	TLQ	*+2	
+17	XCA		
+18	CAS	HOHIGH	
+19	TRA	IHX	ILLEGAL HO BCUNDS
+20	TRA	IHX	
+21	STO	IHA	
+22	STQ	IHB	
+23	LXA	LINEA,1	
+24	TXI	*+1,1,7	
+25	STZ	HONCNT	
+26	CLA	HOTBLE	
IH1	LDQ*	LINEA	
+1	TQP	*+2	
+2	TRA	IH2	NO ENTRIES FOR IMAG. LINES
+3	STA*	LINEA	
+4	ADD	HOLINE	
+5	XCA		
+6	CLA	HONCNT	
+7	ADD	NODES	
+8	STO	HONCNT	
+9	XCA		
IH2	TXI	*+1,1,-1	
+1	TXH	IH1,1,7	
+2	CAS	UPPER	
+3	TRA	IHX	
+4	TRA	IHX	
+5	CAS	LOWER	
+6	TRA	IH3	
+7	TRA	IHX	
+8	TRA	IHX	
IH3	STA	IH7	
+1	SUB	HOTBLE	
+2	PAX	,1	
+3	SXD	HOTBLE,1	

```

+4 CLA IHB
+5 TZE **3
+6 SUB IHA
+7 TNZ IH6
+8 CLA MCNCNT
+9 XCA
+10 MPY IHA
+11 STO HOTOT1
+12 STQ HOTOT
+13 CAL IHA
+14 ALS 9
+15 GRA IHA
+16 ALS 9
+17 OKA IHA
+18 ALS 9
+19 GRA IHA
+20 SLW* IH7
+21 TIX *-1,1,1
+22 ZET IHB
+23 NZT DSPY
+24 TRA **2
+25 TRA IH6
+26 TSX BH,4
      RETURN IH
+27 TRA IH+1
IH6  SUB HOINC
      +1 SLW IHD
      +2 LXA NODES,4
IH4  STZ IHC
      +1 AXT 4,2
IH5  LDQ RND
      +1 MPY K.R
      +2 STQ RND
      +3 MPY IHD
      +4 ADD IHB
      +5 ANA K.M5
      +6 TXL **2,4,0
      +7 TXI **4,4,-1
      +8 CAL HOHIGH
      +9 ORS IHC
+10 TRA IH9
+11 ORS IHC
+12 ADD HCTCT
+13 STO HOTCT
+14 PBT
+15 TRA **4
+16 CLA K.A1
+17 ADD HOTOT1
+18 STO HOTOT1
IH9  CAL IHC
      +1 TNX IH7,2,1
      +2 ALS 9
      +3 SLW IHC
      +4 TRA IH5

```

```

IF HPRIME IS ZERC
CR ENTERED FROM MCDFLS CR BEST-PATH,
USE BEST-PATH ALGORITHM.
ELSE, FILL HCTBLE.
PRESET TO BEST VALUES

```

```
IH7  SLW  **0,1
      +1 TXH  **2,4,0
      +2 LXA  NODES,4
      +3 TIX  IH4,1,1
      RETURN IH
      +4 TRA  IF+1
      RETURN IH,1
IHX  AXT  1,4
```

BH PRESETS HO TABLE TO 'BEST' VALUES

CALL SEQ TSX BH,4
RETURN

	BEGIN	1,7	
BH	TXL	**+5,0,0	SUBROUTINE LINKAGE
+8	CLA	UPDATE	SAVE HO-TABLE
+9	STO	BHX	UPDATING ALGORITHM
+10	CLA	MIN	USE MIN AS
+11	STO	UPDATE	UPDATER
+12	LXA	NODES,4	I = NR. OF NODES,N
BH1	SXA	K.O,4	RECYCLE ON I
+1	LXA	NODES,2	
+2	CLA	HOMAX	
+3	STO*	N2B	HO(K) = MAXHC , K = U(U)N
+4	TIX	*-1,2,1	
+5	STZ*	N2C	HO(I) = C
+6	STZ	BH.A	ACTION-TOGGLE OFF
+7	STL	BH.M	MOVE-TOGGLE ON
+8	TRA	BH3	INTO INNER LOOP WITH J = I
BH2	STZ	BH.A	BEGIN INNER LOOP, ACTION OFF
+1	LXA	NODES,4	J = N
BH4	STZ	BH.M	RECYCLE ON J, MOVE IS OFF
BH3	SXA	K.N,4	
+1	CLA*	N2C	B.HO = HO(J)
+2	STO	B.HO	K = C, (COUNT FROM 8-K TO 1)
+3	AXT	8,2	
+4	PXA	,4	8*J+K IS REL. PCS. OF LINE(J,K)
+5	ALS	3	IN LINE TABLE
+6	PAX	,4	
+7	SXD	BH9,4	LINE(J,K)
BH5	CLA*	BUSYC	IGNORE IF LINE DEAD
+1	TMI	BH6	
+2	PDX	,1	IGNORE IF NO MSG ON LINE
+3	TXL	BH6,1,0	ERASE MSG
+4	PXA		
+5	STD*	BUSYC	
+6	PXA	,1	FOR FUTURE USE BY HOVER
+7	STO	K.HO	B.HO = MIN(B.HC, HC OF MSG)
+8	CAS	B.HC	
+9	TRA	**+4	
+10	TRA	**+3	
+11	STO	B.HO	
+12	STL	BH.M	MOVE IS ON IF K.HC DECREASES
+13	SXA	K.L,4	SAVE LINE NR
+14	XEC	HOVER	UPDATE HO-TABLE ENTRY
+15	LXA	K.L,4	
BH6	TXN	**+2,2,1	INCREMENT K, TEST
+1	TXI	BH5,4,1	AND RECYCLE IF K LESS THAN 8
+2	N2T	BH.M	TEST MOVE TOGGLE
+3	TRA	BH10	OFF, NODE J SENDS NO MSGS
+4	CLA	B.HO	ON

```

+5 LXA K.N,2
+6 STO* N2B
+7 ADD HCINC
+8 ALS 18
+9 STO K.HO
BH8 CLA* BUSYC
+1 TMI BH9
+2 PAX ,2
+3 CLA K.HO
+4 STD* BUSYB
BH9 TXL **2,4,**0
+1 TXI BH8,4,-1
+2 STL BH.A
BH10 LXA K.N,4
+1 TIX BH4,4,1
+2 ZET BH.A
+3 TRA BH2
+4 LXA K.O,4
+5 TIX BH1,4,1
+6 LXA NODES,4
+7 STZ* N2C
+8 TIX *-1,4,1
+9 CLA BHX
+10 STO UPDATE
      RETURN BH
+11 TRA BH+1
BHX PZE

```

```

      RESET HC(J) AND INC. B.HC --
      SENT TO ALL NEIGHBORS

```

```

      IGNORE DEAD LINES

```

```

      IMAGE LINE
      MSGS SENT
      TO NEIGHBORS
      FINI IF K = 0
      DECREMENT K AND RECYCLE

```

```

      RECYCLE INNER LCCP ON J

```

```

      TEST ACTION TOGGLE
      ON, REDO INNER LCCP
      OFF
      RECYCLE ON I
      FINI, CLEAR THINGS

```

```

      RESTORE UPDATING
      ALGORITHM

```

MCVER MOVES MESSAGES THRU NET, UPDATES HO TABLE
FOR REROUTED MSGS., STACKS REROUTED MSGS. AT
APPROPRIATE NODE, PROCESSES DELIVERED MSGS.

CALL SEQ XEC MCVER
NORMAL RETURN
SAVES IR1, IR2 AND SENS

MA IS A MOVER

MA	TSX	*+1,4	
	BEGIN	1,7,1	
MA0	TXL	*+7,0,0	SUBROUTINE LINKAGE
+11	LXA	TP,1	
+12	TXI	*+1,1,1	NEXT FRAME'S TRAFFIC LIST
+13	TIX	*+1,1,TPU	COLNT MODULC TPL
+14	SXA	TP,1	
+15	CLA	TP,1	
+16	STZ	TP,1	
+17	PDX	,2	
+18	TRA	MA1+1	
MA1	LXD	TP,2	
+1	TXH	MA3,2,0	
	RETURN	MA0	END TRAFFIC
+2	TRA	MA0+1	
MA3	CLA	MSG2,2	
+1	STC	TP	
+2	SXA	K.MSG,2	
+3	PAX	,1	EFFERENT LINE NR.
+4	CLA*	BUSYA	
+5	PAX	,1	IMAGE LINE
+6	STA	K.L	
+7	ARS	3	
+8	PAX	,4	
+9	STA	K.N	AFFERENT NODE
+10	CLA*	NIC	DROP MSG
+11	TMI	MA2.1	
+12	LDQ	MSG1,2	
+13	PXA		
+14	LGL	7	
+15	STO	K.O	ORIGIN
+16	PXA		
+17	LGL	7	
+18	STC	K.D	DESTINATION
+19	PXA		
+20	LGL	4	
+21	STO	K.P	PRIORITY
+22	PXA		
+23	LGL	9	
+24	STC	K.ST	TIME DELAYED INSTACKS
+25	PXA		
+26	LGL	9	
+27	ADD	HOINC	HO
+28	STO	K.HO	
+29	XEC	HOVER	UPDATE HC TABLE
+30	CLA	K.N	
+31	SUB	K.D	
+32	TNZ	MA4	
+33	XEC	RECORD	RECORD ARRIVAL
+34	TRA	MA1	
MA4	CLA	K.HO	
+1	CAS	HCMAX	DROP MSG IF HO EXCESSIVE

	+2	TRA	MA2
	+3	NOP	
	+4	CAL	MSG1,2
	+5	ADD	HGINCI
	+6	SLW	MSG1,2
	+7	XEC	DELIVR
	+8	TRA	MA1
MA2		CAS	HCHI1
	+1	TRA	**4
	+2	TRA	**3
MA2.1		XEC	DROP
	+1	TRA	MA1
	+2	CLA	HCHI1
	+3	STO	K.HO
	+4	TRA	MA1

INC HO
IN MSG TABLE
STACK MSG

HOVER UPDATES HOTBLE FOR LINE I, MESSAGE J
AFTER INCREMENTING HO IN MSG TABLE

```
CALL SEQ. XEC HOVER
      RETURN
      (K.L) = I
      (K.MSG) = J
      (K.N) = CURRENT NODE
      (K.O) = ORIGIN(J)
      (K.HO) = HO(J)
      (K.D) = DEST(J)
      (K.P) = PRIORITY(J)
      SAVES IR1, IR2 AND SENS
```

HA IS A HOVER

HA	TSX	**1,4	
	BEGIN	1,7,1	
HA0	TXL	**7,0,C	SUBROUTINE LINKAGE
+11	CLA	K.O	
+12	CAS	K.N	
+13	TRA	**2	
+14	STZ	K.HO	
+15	SUB	K.A1	
+16	LGR	2	
+17	PAC	,4	REL WRD NR CF HC-ITEM WRT CRIGIN
+18	PXA		
+19	LGL	2	
+20	PAX	,2	POSITION IN WORD CF ITEM
+21	LXA	K.L,1	
+22	CLA*	LINEA	
+23	STA	HA4	
HA4	CAL	**0,4	HO WORD
+1	XEC	HSH,2	SHIFT
+2	ANA	K.M7	
+3	LDQ	K.HO	
+4	SXA	HA1,4	
+5	XEC	UPDATE	EXECUTE UPDATING ALGORITHM
+6	TRA	HA2	
HA1	AXT	**0,4	
+1	STO	HA.1	
+2	CAL	K.M7	
+3	XEC	HSH1,2	
+4	COM		
+5	ANA*	HA4	
+6	STQ*	HA4	
+7	ORS*	HA4	
+8	CLS	HA.1	
+9	ADD	K.HO	
+10	ADD	HOTOT	
+11	SLW	HOTOT	
+12	TZE	HA5	
+13	TMI	HA3	
+14	PBT		
+15	TRA	HA5	
+16	CLA	HOTOT1	
+17	ADD	K.A1	
+18	STO	HOTOT1	
HA5	EQU	HA2	
	RETURN	HA0	
HA2	TRA	HA0+1	
HA3	CLA	HOTOT1	
+1	SUB	K.A1	
+2	SLW	HOTOT1	
	RETURN	HA0	
+3	TRA	HA0+1	

UPDATE EXECUTES HO-TABLE UPDATING ALGORITHM.
 CALL SEQ XEC UPDATE
 RETURN IF NO CHANGE IN HO-TABLE
 RETURN IF TABLE TO BE UPDATED
 ENTRY (ACC) = OLD ENTRY
 (MQ) = HO OF MESSAGE
 EXIT (MQ) = NEW ENTRY
 SAVES IR'S AND SENSE

MIN UPDATES BY MAKING NEW ENTRY
 = MIN(OLD ENTRY, HO OF MSG)

MIN TLQ HA1

LRN UPDATES HO ENTRIES AS FOLLOWS --
 $HO = HO + K1 * (MSG\ HO - HO)$ IF (MSG HO - HO) IS NEG.
 $HO = HO + K2 * (MSG\ HO - HO)$ IF (MSG HO - HO) IS POS.

LRN	TSX	LRNO,4	HO ENTRY
LRNO	STO	LRN1	
+1	XCA		
+2	SUB	LRN1	MSG HO - HO
+3	TZE	1,4	NO CHANGE
+4	XCA		
+5	TQP	**3	
+6	MPY	K1	LEARNING - USE K1
+7	TRA	**2	
+8	MPY	K2	FORGETTING - USE K2
+9	ADD	LRN1	
+10	XCA		
+11	CLA	LRN1	
+12	TRA	2,4	

DELIVR APPENDS MESSAGE J TO TRAFFIC STACK
FOR CURRENT NODE

CALL SEQU XEC DELIVR
RETURN
INITIAL CONDITIONS AS FOR HOVER
SAVES IR1, IR2 AND SENS

DA IS A DELIVR

DA STACKS MESSAGES AS FOLLOWS --

DECREMENT (ADDRESS) OF NODE1(J) POINTS TO FIRST
ENROUTE (NEW) MSG ON STACK FOR NODE J.

FOR ALL MSGS IN STACK, HD IS INCREMENTED AND MSG2 IS

PREFIX = MZE

DECR = POINTER TO NEXT MSG

TAG = INCOMING LINK NUMBER

ADDRESS = TIME IN STACK (ORIGINALLY ZERO)

A ZERO LINE NR (K.L) INDICATES NEW MSG.

A ZERO POINTER INDICATES ENC OF STACK

DA	TSX	DA0,4	
	BEGIN	1,7,1	
DA0	TXL	**7,0,0	SUBROUTINE LINKAGE
+11	LXA	K.MSG,2	
DA1	CAL	K.L	
+1	ALS	15	
+2	ANA	K.M3	STACK-FLAG ON, LINK NR IN TAG
+3	ORA	K.MZE	
+4	SLW	MSG2,2	
+5	LXA	K.N,1	
+6	CLA*	N1A	
+7	PDX	,4	
+8	ZET	K.L	TEST MSG
+9	TRA	DA4	ENROUTE MSG
+10	NZT	CHOKE	NEW MSG. ARE WE CHOKING INPUT
+11	TRA	DA4	NO - TREAT AS ENROUTE MSG.
+12	PAX	,4	YES - PLACE ON SPECIAL INPUT STACK
+13	TXH	DA2,4,0	IS STACK EMPTY - NO
+14	PXA	,2	YES
+15	STA*	N1A	
	RETURN	DA0	
+16	TRA	DA0+1	
DA4	TXH	DA2,4,0	IS STACK EMPTY - NO
+1	PXD	,2	YES
+2	STD*	N1A	
	RETURN	DA0	
+3	TRA	DA0+1	
DA2	SXA	DA3,4	
+1	CLA	MSG2,4	
+2	PDX	,4	
+3	TXH	DA2,4,0	
DA3	AXT	**0,4	PUT MSG ON END OF STACK
+1	PXD	,2	
+2	STD	MSG2,4	
	RETURN	DA0	
+3	TRA	DA0+1	

RA IS A MESSAGE STACK PROCESSOR.
 FOR EACH MESSAGE ON STACK FOR NODE J, RA CHOOSES
 NON-BUSY EFFERENT LINE HAVING SMALLEST HANDCOVER
 NUMBER TO MESSAGE'S DESTINATION. TWO STACKS ARE
 USED - A 'STANDARD' STACK AND A 'NEW-MESSAGE'
 STACK, THE LATTER BEING USED FOR INPLT CHCKING ONLY.
 IF NO LINES ARE AVAILABLE, MSG REMAINS CN STACK. THE
 STANDARD STACK IS PROCESSED BEFORE NEW-MESSAGE STACK,
 AND MESSAGES IN LATTER STACK ARE NEVER DROPPED -
 THUS STACKING AT INPUT IS ENFORCED.

CALL SEQUENCE XEC RA
 RETURN
 SAVES IR'S AND SENSE

RA	TSX	RA0,4	
	BEGIN	1,7,1	
RA0	TXL	++7,0,0	SUBROUTINE LINKAGE
+11	LXA	LINES,1	
+12	TXI	++1,1,7	
RA03	CLA*	BUSYA	FOR ALL LINES, ---
+1	TMI	RA02	IGNORE DEAD LINES
+2	PDX	,2	
+3	TXL	RA02,2,0	IGNORE FREE LINES
+4	TXI	++1,2,-1	DECREMENT
+5	PXD	,2	BUSY
+6	STD*	BUSYA	COUNTER
+7	TNZ	RA02	
+8	CLA*	LINEA	FREE LINE IF COUNT BECCMES ZERO
+9	SLW*	LINEA	
RA02	TIX	RA03,1,1	
+1	LXA	NODES,2	
RA00	SXA	K.N,2	
+1	STZ	RA.0	FIRST CYCLE PROCESSES ENROUTE MSGS
+2	STZ	RA.4	BUSY-FLAG OFF
+3	STZ	RA.5	STACK-USE FLAG CFF
+4	STZ*	N2B	CLEAR STACK COUNT
+5	CLA*	N1B	
+6	STO	MSG2	
RA01	AXT	,1	
RA1	SXA	RA.1,1	LAST MSG NR.
+1	CLA	MSG2,1	
+2	PDX	,1	
+3	TXH	RA3,1,0	TEST END-CF-STACK
RA2	LDQ	MSG2	END
+1	RQL	1B	
+2	STQ	MSG2	
+3	ZET	RA.0	TEST CYCLE
+4	TRA	RAX	LAST - FINI
+5	STL	RA.0	
+6	TRA	RA01	
RAX	STQ*	N1B	

	+1	TIX	RA00,2,1	
		RETURN	RA0	
	+2	TRA	RA0+1	
RA3	+1	SXA	K.MSG,1	CURRENT MSG NR.
	+1	ZET	RA.4	ARE ALL LINES BUSY
	+2	TRA	*+4	YES
	+3	TSX	FA,4	ORDER AVAIL. LINES FOR ROUTING
	+4	LXA	K.AL,4	NR. OF AVAILABLE LINES
	+5	TXH	RA4,4,0	ANY AVAIL. LINES
	+6	STL	RA.4	NO, SET BUSY FLAG
	+7	ZET	RA.0	TEST CYCLE
	+8	TRA	RA9	SECOND CYCLE (INPUT STACK)
RA6		CLA*	N2B	FIRST CYCLE (STANDARD STACK)
	+1	CAS	STACK	IS STACK FULL
	+2	TRA	RA7	YES
	+3	TRA	RA7	YES
	+4	ADD	K.A1	NO, INCREMENT COUNT
	+5	STO*	N2B	
	+6	CAL	MSG2,1	INCREMENT TIME-IN-STACK
	+7	ADC	K.A1	
	+8	STA	MSG2,1	
	+9	STL	RA.5	STACK IN USE
RA9		CAL	MSG1,1	INC STACK-TIME COUNTER
	+1	ADD	STKINC	
	+2	SLW	MSG1,1	
	+3	TRA	RA1	
RA7		CLA	MSG2,1	
	+1	LXA	RA.1,1	DELETE MSG FROM STACK
	+2	STD	MSG2,1	
	+3	XEC	DROP	
	+4	TRA	RA1	
RA4		CLA	AL,4	ROUTE MSG OVER THIS LINE
	+1	PAX	,4	
	+2	STA	MSG2,1	LINE NR. TO MSG TABLE
	+3	STP	MSG2,1	STACK-FLAG OFF
	+4	CAL	K.MZE	MAKE LINE BUSY
	+5	STP*	LINEC	
	+6	CLA	GRAIN	
	+7	ADD	K.A1	
	+8	ALS	18	
	+9	ADD*	BUSYC	MAKE LINE BUSIER BY GRAIN+1
	+10	STO*	BUSYC	
	+11	STD	RA8	
	+12	CLA	MSG2,1	
	+13	LXA	RA.1,1	DELETE MSG FROM STACK
	+14	STD	MSG2,1	
	+15	CLA*	LINEC	SET MSG TP
	+16	STD	*+3	MSG ARRIVES NEXT STATION AT TIME
	+17	LXA	TP,1	T* + TP(LINE) + BUSY(LINE)
RA8		TXI	*+1,1,**0	+ GRAIN + 1
	+1	TXI	*+1,1,**0	
	+2	TIX	*+1,1,TPU	MODULO TPU
	+3	TSX	CHAIN,4	STACK MSG
	+4		TP,1	ON LIST FOR TIME-FRAME-NR

+5	LXA	RA.1,1
+6	NZT	SNAP
+7	TRA	RA1
+8	CLA*	N4B
+9	ADD	K.A1
+10	STO*	N4B
+11	TRA	RA1

COUNT MESSAGES PASSING
THRU NCDE IF SNAPSHCTS CF FLOW
ARE BEING TAKEN.

FA COMPILES, FOR GIVEN NODE AND MSG, A LIST OF
AVAILABLE AFF. LINES ORDERED BY INCREASING H.C. NR.
W.R.T. DESTINATION NCDE OF MSG.
LINES HAVING IDENTICAL H.O. NRS. ARE ORDERED
RANDOMLY WITHIN THEIR COMMON H.O. NR.

CALL SEQ. TSX FA,4
RETURN
(IR1) = MESSAGE NR. = (K.MSG)
(IR2) = NODE NR. = (K.N)
EXIT (K.AL) = NR. AVAILABLE LINES
(AL(I)) = I-TH BEST CHOICE OF LINES
ADD.(AL(I)) = LINE NR.
BITS 0-9 = C
BITS 10-21 = ASSOC. HO NR FOR DEST. NCDE

	BEGIN	1,7	
FA	TXL	**+5,0,0	SUBROUTINE LINKAGE
+8	STZ	K.AL	
+9	PXA	,2	
+10	ALS	3	
+11	PAX	,2	FIRST LINE NR. FROM SOURCE NODE
+12	LDQ	MSG1,1	
+13	LGL	7	
+14	PXA		
+15	LGL	7	DESTINATION NODE
+16	STO	K.D	
+17	SUB	K.A1	
+18	LGR	2	
+19	STA	FA2	REL. HO WORD WRT TC DEST NODE
+20	PXA		
+21	LGL	2	
+22	PAX	,4	POSITION IN WORD OF ITEM
+23	CLA	HSH2,4	SHIFT
+24	STO	FA3	
+25	AXT	8,1	
FA1	CLA*	LINEB	RECYCLE IF LINE BUSY
+1	TMI	FA6	
+2	SXA	K.L,2	
+3	STA	**+2	
FA2	AXC	**0,4	
+1	CAL	**0,4	FETCH HO WORD
FA3	NOP		SHIFT
+1	ANA	K.M1	
+2	STO	ALHO	
+3	STO	ALWD	
+4	SXA	ALWD,2	SAVE AL ENTRY PZE LINE,,H.O.
+5	TSX	FB,4	INSERT NEW AL ENTRY
FA6	TXI	**+1,2,1	
+1	TIX	FA1,1,1	
	RETURN	FA	
+2	TRA	FA+1	

FB UPDATES AL LIST, INSERTING ENTRY
 FOUND IN (ALWD), ORDERED BY H.C. NR. IN (ALFC).
 AL COUNT IN (K.AL)

	BEGIN	1,7	
FB	TXL	**5,0,0	SUBROUTINE LINKAGE
+8	LXA	K.AL,1	
+9	TXH	**2,1,0	
+10	TXI	FB2,1,1	
+11	LDQ	RND	
+12	MPY	K.R	
+13	STC	RND	
+14	RQL	12	
FB1	CAL	AL,1	
+1	ANA	K.M1	
+2	CAS	ALHO	
+3	TXI	FB2,1,1	INSERT ABOVE HERE
+4	TRA	FB3	EQUALITY - RESOLVE CONFLICT
FB6	TIX	FB1,1,1	
FB2	SXC	FB4,1	SAVE PCINTER
+1	LXA	K.AL,1	
+2	TXI	**1,1,1	INCREMENT COUNTER
+3	SXA	K.AL,1	
FB4	TXL	FB5,1,**C	
+1	CLA	AL+1,1	
+2	STO	AL,1	MOVE UP
+3	TXI	FB4,1,-1	
FB5	CLA	ALWD	INSERT HERE
+1	STO	AL,1	
	RETURN	FB	
+2	TRA	FB+1	
FB3	RQL	1	
+1	TQP	FB6	FLIP COIN TO RESOLVE CONFLICT
+2	TXI	FB2,1,1	

DROP APPENDS DISCARDED TRAFFIC TO DRCP-LIST.
 USING MSG TABLE AS STACK. MSG2 BECOMES
 MZE CURRENT NODE , X , POINTER TO NEXT MSG
 WHERE X = 4 IF MSG DRCPED IN STACK, ELSE 0.
 TOTAL DROPS AND DROPS AT NODE ARE INCREMENTED.

CALL SEQU XEC DROP
 RETURN

(K.N) = NOCE NR
 (K.MSG) = MSG NR
 SAVES IR1, IR2 AND SENSE

DROP	TSX	**+1,4	
	BEGIN	1,7	
DRO	TXL	**+5,0,0	SUBROUTINE LINKAGE
+8	LXA	K.MSG,1	
+9	CAL	MSG2,1	
+10	AXT	,2	
+11	PBT		WAS MSG DRCPED INSTACK
+12	AXT	1,2	NO
+13	ARS	18	
+14	ANA	K.M3	TAG = 4 IF IN STACK
+15	SSM		PREFIX = MZE
+16	ORA	K.N	ADD = NODE
+17	STO	MSG2,1	
+18	CLA	STKCNT,2	
+19	ADD	K.A1	COUNT DROPS
+20	STO	STKCNT,2	
+21	CLA	MSGCNT	
+22	SUB	K.A1	
+23	STO	MSGCNT	
+24	TSX	CHAIN,4	
+25		DR	
	RETURN	DRO	
+26	TRA	DRO+1	

RECORD APPENDS COMPLETED TRAFFIC TO ARRIVAL-LIST

CALL SEQ XEC RECCRD
 RETURN
 (K.MSG) = MSG NR.
 SAVES IRI, IR2 AND SENSE

RECORD	TSX	#+1,4	
	BEGIN	1,7	
RCO	TXL	#+5,0,0	SLBROUTINE LINKAGE
+8	LXA	K.MSG,2	
+9	CLA	K.MZE	
+10	STO	MSG2,2	
+11	CLA	ARRCNT	
+12	ADD	K.A1	COUNT ARRIVALS
+13	STO	ARRCNT	
+14	CLA	K.HO	
+15	ADD	ARRHO	
+16	STO	ARRHO	
+17	PBT		
+18	TRA	#+4	
+19	CLA	ARRHO1	
+20	ADD	K.A1	
+21	STO	ARRHO1	
+22	CLA	K.HC	
+23	TZE	#+6	
+24	AKS	3	
+25	PAX	,1	
+26	CLA*	HODA	
+27	ADD	K.A1	
+28	STO*	HODA	
+29	CLA	K.ST	
+30	PAX	,1	
+31	TXI	#+1,1,1	
+32	CAS	MAXHO	
+33	NOP		
+34	LXA	MAXHO,1	
+35	CLA*	STDA	
+36	ADD	K.A1	
+37	STO*	STDA	
+38	CLA	MSGCNT	
+39	SUB	K.A1	
+40	STO	MSGCNT	
+41	TSX	CHAIN,4	
+42		AR	
	RETURN	RCO	
+43	TRA	RCO+1	

GN GENERATES A RANDOM NODE NR

```
CALL SEQ   TSX   GN,4
          RETURN
EXIT      (ACC) = NODE NR
          SAVES IR1, IR2 AND SENSE
```

```
GN      BEGIN   1,4
          TXL    **3,0,0
          LDQ    RND
+4      LDQ    RND
+5      MPY    K.R
+6      STQ    RND
+7      MPY    NODES
+8      ACC    K.A1
          RETURN GN
+9      TRA    GN+1
```

SUBROUTINE LINKAGE

CHAIN ADDS MSG. TO TOP OF LIST CHAINED VIA DECR.
 HEAD OF LIST IS PZE LAST MSG NR , , FIRST MSG NR

CALL SEQ

TSX CHAIN,4
 PZE LIST-NAME (MAY BE TAGGED)
 RETURN
 (K.MSG) = MSG NR
 SAVES IR1, IR2 AND SENS

CHAIN	BEGIN	2,4	
	TXL	#+3,0,0	SUBROUTINE LINKAGE
+4	LDQ*	1,4	
+5	CLA	K.MSG	
+6	STA*	1,4	UPDATE LIST'S TAIL
+7	ALS	18	
+8	XCA		
+9	TNZ	CH1	
+10	XCA		EMPTY LIST
+11	STD*	1,4	HEAD = TAIL
+12	TRA	CH3	
CH1	PAX	,4	
+1	XCA		
+2	STD	MSG2,4	CHAIN LAST TO CURRENT
CH3	PCX	,4	
+1	PXA		
+2	STD	MSG2,4	CURRENT BECOMES LAST
	RETURN	CHAIN	
CH2	TRA	CHAIN+1	

MG PRESETS MSG STACKS AND GENERATES DESIRED
NR. OF MSGS.

	CALL	SEQ	TSX	MG,4	
			RETURN		
	BEGIN	1,7			
MG	TXL	**5,0,0			SUBROUTINE LINKAGE
	+8 CLA	MSG5			
	+9 ALS	1			
	+10 PAC	,1			
	+11 STZ	MSGCNT			
MG1	SXA	K.MSG,1			
	+1 TSX	GM,4			
	+2 NOP				SYSTEM LOADED
	+3 TXI	**1,1,2			
	+4 TXH	MG1,1,0			
	+5 STZ	AR			
	+6 STZ	DR			
	+7 STZ	OL			
	RETURN	MG			
	+8 TRA	MG+1			

ACTIVE RESETS ACTIVITY LEVELS
 CALL SEQU TSX ACTIVE,4
 RETURN

	BEGIN	1,7	
ACTIVE	TXL	**5,0,0	SLBROUTINE LINKAGE
+8	LDQ	IMPARE	
+9	FMP	FACTOR	
+10	CAS	MAXACT	IS DESIRED ACTIVITY TCC LARGE
+11	TRA	**3	YES
+12	TRA	**2	YES
+13	TRA	**5	NO
+14	CLA	MAXACT	USE MAXIMUM
+15	STO	ALPHA	
+16	CLA	MSGs	
+17	TRA	ACTV1	
+18	STO	ALPHA	
+19	FDP	MAXACT	
+20	PXA		
+21	LLS	8	
+22	SUB	KINT2	
+23	STA	**2	
+24	PXA		
+25	LRS	**0	
+26	MPY	MSGs	
ACTV1	STO	LOAD	
+1	NZT	PRETGL	
	RETURN	ACTIVE	
+2	TRA	ACTIVE+1	
	XEJECT		
+3	STL	SYSOED	
	XPRINT	D,2	
+6	STL	SYSOED	
+9	PZE	PM-RANDCM+ALPHA,,1	
+10	PZE	ALPHA,,1	
	RETURN	ACTIVE	
+11	TRA	ACTIVE+1	

GM GENERATES A RANDOM MSG WHOSE MSG. NR. IS
(K.MSG). IF DESIRED MSG LOADING HAS NOT BEEN
ATTAINED, MSG ENTERS SYSTEM VIA DELIVR -
OTHERWISE MSG IS PLACED ON DROP LIST

CALL SEQU TSX GM,4
RETURNED IF LOADED
NORMAL RETURN

	BEGIN	2,7	
GM	TXL	*+5,0,0	SLBRCLTINE LINKAGE
GM1	LXA	SOURCE,1	GENERATE CRIGIN
	+1	TXL GM3,1,0	NO SCURCE DIST.
	+2	LXA SCH,4	
	+3	LDQ RND	USE SOURCE DIST.
	+4	MPY K.R	
	+5	STQ RND	
	+6	XCA	
GM4	TXI	*+1,4,1	
	+1	CAS* SC1A	
	+2	TRA GM3	
	+3	NOP	
	+4	CAS* SC2A	
	+5	TRA GM2	
	+6	TRA GM2	
	+7	TIX GM4,1,1	
GM3	LDQ	RND	USE UNIFORM DIST.
	+1	MPY K.R	
	+2	STQ RND	
	+3	MPY SCH	
	+4	ADC K.A1	
	+5	PAX ,4	
GM2	CLA*	N3C	
	+1	STZ K.N	
	+2	STA K.N	
	+3	PAX ,4	
	+4	CLA* NIC	
	+5	TMI GM1	DISCARD IF DEAD
GM5	LXA	SINK,1	GENERATE DESTINATION
	+1	TXL GM7,1,0	NO SINK DIST.
	+2	LXA SKH,4	
	+3	LDQ RND	USE SINK DIST.
	+4	MPY K.R	
	+5	STQ RND	
	+6	XCA	
GM8	TXI	*+1,4,1	
	+1	CAS* SK1A	
	+2	TRA GM7	
	+3	NOP	
	+4	CAS* SK2A	
	+5	TRA GM6	
	+6	TRA GM6	

+7 TIX GM8,1,1
 GM7 LDQ RND
 +1 MPY K.R
 +2 STQ RND
 +3 MPY SKH
 +4 ADD K.A1
 +5 PAX ,4
 GM6 CLA* N3C
 +1 ARS 18
 +2 CAS K.N
 +3 TRA **2
 +4 TRA GM5
 +5 STO K.D
 +6 ALS 29
 +7 XCL
 +8 CLA K.N
 +9 LGR 7
 +10 LXA K.MSG,2
 +11 STQ MSG1,2
 +12 STZ MSG2,2
 +13 CLA MSGCNT
 +14 CAS LCAD
 +15 TRA GM9
 +16 TRA GM9
 +17 ADD K.A1
 +18 STO MSGCNT
 +19 STZ K.L
 +20 XEC DELIVR
 RETURN GM
 +21 TRA GM+1
 GM9 TSX CHAIN,4
 +1 PZE DR
 RETURN GM,1
 +2 AXT 1,4

DISCARD IF INCESTVCUS

ORG, HO=C, DEST, PRICRITY=C

IS SYSTEM LOADED

YES

YES

NO

INITIALIZE FCR DELIVR
 STACK MESSAGE

MSG TO DROP LIST

RM REPLACES DROPPED AND DELIVERED MSGS
WITH RANDOMLY INITIATED FRESH MSGS.

CALL SEQU TSX RM,4

	BEGIN	1,7	
RM	TXL	**5,0,0	SUBROUTINE LINKAGE
RM1	LXD	AR,1	
	+1	TXL RM2,1,0	
	+2	CLA MSG2,1	
	+3	STD AR	
	+4	SXA K.MSG,1	
	+5	TSX GM,4	
		RETURN RM	SYSTEM LOADED
	+6	TRA RM+1	
	+7	TRA RM1	
RM2	LXD	DR,1	
	+1	TXL RM3,1,0	
	+2	CLA MSG2,1	
	+3	STD DR	
	+4	SXA K.MSG,1	
	+5	TSX GM,4	
		RETURN RM	SYSTEM LOADED
	+6	TRA RM+1	
	+7	TRA RM2	
RM3	STZ	AR	
	+1	STZ DR	
	+2	TRA RM+1	

SIMUL8 CYCLES THRU SIMULATOR K TIMES AFTER
INITIALIZING NET.

CALL SEQU TSX SIMUL8,4
PZE K OR LOC OF K
PZE L,TAG
VFD H36/K OR LOC OF K
RETURN IF NOGC
NORMAL RETURN

WHERE (L,TAG) CONTAINS -
LOC OF FIRST SPECIAL LINK,,NR. OF SPECIAL LINKS
AND L=TAG=0 IMPLIES NO SPECIAL LINKS.

	BEGIN	5,7,1	
SIMUL8	TXL	**7,0,0	SUBROUTINE LINKAGE
+11	STZ	ACTTGL	
+12	CLA	1,4	
+13	STO	SIM1	
+14	CLA	3,4	
+15	STO	SIM1+1	
+16	CLA	2,4	
+17	TZE	**2	
+18	CLA*	2,4	
+19	STO	**2	
+20	TSX	PRES,4	INITIALIZE NET, DISPLAY PARAMS
+21	PZE		
+22	TRA	SIM2	ERROR
+23	STZ	INITGL	
+24	TSX	CONT,4	
SIM1	PZE		
	VFD	H36/	
+1			
	RETURN	SIMUL8	
+2	TRA	SIMUL8+1	
	RETURN	SIMUL8,1	
SIM2	AXT	1,4	

CONTINUE SIMULATION FOR K MORE CYCLES

CALL SEQU TSX CONT,4
PZE K OR LOC OF K
VFD H36/K OR LOC OF K

	SPACE	2	
	BEGIN	3,7,1	
CONT	TXL	**7,0,0	SUBROUTINE LINKAGE
NGO	EQU	CONT	
-2	AXT	6,1	TEST IF K OR LCC CF K.
-1	LDQ	2,4	
NGC	PXA		
+1	LGL	6	

	+2	CAS	K.A48	TEST IF BLANK--
	+3	TRA	NGB	IMPLIES LOCATICN
	+4	TRA	**4	BLANK
	+5	CAS	K.A9	TEST IF NUMERIC
	+6	TRA	NGB	IMPLIES LOCATION
	+7	NOP		NUMERIC--
	+8	TIX	NGC,1,1	DITTO.
	+9	CLA	1,4	NUMERIC IMPLIES K
	+10	TRA	**2	
NGB		CLA*	1,4	
	+1	SSP		
	+2	TZE	NGO+1	EXIT IF ZERO CYCLE COUNT.
	+3	CAS	K.M2	TEST IF FLOATING
	+4	TRA	**3	
	+5	TRA	NGD+1	
	+6	TRA	NGD+1	
	+7	XCA		CONVERT
	+8	PXA		TO
	+9	LLS	8	INTEGER
	+10	SUB	KINT2	
	+11	TPL	**2	
	+12	PXA		
	+13	STA	NGD	
	+14	PXA		
NGD		LLS	**0	
	+1	NOP		
				CALCULATE OUTPUT INTERUPT FREQUENCY FROM K
	+2	STO	NG3	
	+3	LDQ	NG3	
	+4	ARS	1	
	+5	TZE	**5	
	+6	STO	NG3	
	+7	ARS	1	
	+8	TZE	**2	
	+9	STO	NG3	
	+10	XCA		
	+11	ZET	INITGL	IS ENTRY VIA SIMUL8
	+12	TRA	NGA	NO - CONTINUATION
	+13	STO	NG1	YES - CLEAR THINGS
	+14	TSX	CLEAR,4	
	+15	STZ	COUNTS	
	+16	STL	INITGL	
	+17	TSX	MG,4	PRESET MESSAGES - STACK THEM AT NODES
	+18	TRA	**3	
NGA		ADD	NG1	
	+1	STO	NG1	
	+2	CLA	NG3	
	+3	TZE	NGO+1	
	+4	ADD	COUNTS	
	+5	STO	NG4	
	+6	CLA	HMAX	RESET HMAX IN CASE CHANGED
	+7	ALS	3	
	+8	STO	HOMAX	
	+9	LDQ	FACTOR	

	+10	FMP	IMPARE	
	+11	CAS	ALPHA	HASLOADING BEEN CHANGED
	+12	TRA	**2	YES
	+13	TRA	**2	NC
	+14	TSX	ACTIVE,4	YES - CALCULATE NEW ACTIVITY LEVELS
	+15	CLA	LEARN	RESET LEARN AND FORGET
	+16	TSX	FRACT,4	
	+17	LDQ	KFULL	
	+18	STQ	K1	
	+19	CLA	FORGET	
	+20	TSX	FRACT,4	
	+21	LDQ	KFULL	
	+22	STQ	K2	
	+23	LDQ	MIN	SET HO-ENTRY
	+24	ZET	ADAPT	UPDATING ALGRITHM
	+25	LDQ	LRN	
	+26	STQ	UPDATE	
NG2		SWT	5	EXCESS TIME TEST
	+1	TRA	**5	OKAY - CONTINUE
	+2	TSX	SIDIST,4	
	+3	TSX	SIFLOW,4	
	+4	TSX	SIWAIT,4	
	+5	TRA	SYSTEM	
	+6	XEC	ROUTER	ROUTE MSGS AT ACDE
	+7	XEC	MCOVER	MOVE MSGS THRU NETWORK
	+8	TSX	RM,4	REPLACE DLVD/DROPPED MSGS WITH FRESH
	+9	CLA	COUNTS	INC. AND TEST CYCLE CCUNT
	+10	ADD	K.A1	
	+11	STO	COUNTS	
	+12	CAS	NG1	
	+13	TRA	NG5	
	+14	NOP		
	+15	NZT	BFTST	
	+16	TRA	NG2	
	+17	CAS	NG4	
	+18	TRA	**3	
	+19	TRA	**2	
	+20	TRA	NG2	
	+21	ADD	NG3	
	+22	STO	NG4	
	+23	TSX	SYSTST,4	
	+24	TRA	NG2	
NG5		CLA	HOTOT1	
	+1	LDQ	HOTOT	
	+2	LLS	7	
	+3	DVP	HONCNT	
	+4	STQ	HOAVG	MEAN OF HO TABLE
	+5	CLA	DRPCNT	
	+6	ADD	STKCNT	
	+7	STO	K.T1	
	+8	ADD	ARRCNT	
	+9	LDQ	K.T1	
	+10	STO	K.T1	
	+11	MPY	K.A100	

+12 LLS 10
+13 DVP K.T1
+14 STQ DROPPC
+15 CLA ARRHO1
+16 LDQ ARRHO
+17 LLS 7
+18 DVP ARRCNT
+19 STQ ARRAYG
RETURN NGO
+20 TRA NGO+1
NG1 OCT 0

PERCENT MSGS DRCPED

MEAN HC CF DELIVERED MSGS.

SKTLU APPENDS NEW SINK TO SINK TABLE

CALL SEQ TSX SKTLU,4
 PZE NCDE NR
 DEC PERCENT MSGS TO BE PREFEMPTED

SKTLL	BEGIN	3,7	
	TXL	*+5,0,0	SUBROUTINE LINKAGE
+8	CLA	1,4	
+9	STO	SC1	
+10	TNZ	*+2	
+11	TSX	ERROR,4	
+12	CAS	NCDES	
+13	TSX	ERROR,4	
+14	NOP		
+15	CLA	2,4	
+16	TSX	FRACT,4	
+17	TSX	ERROR,4	
+18	STQ	SC2	
+19	LXA	SKH,1	
+20	CLA	SC1	
+21	PAX	,4	
+22	CAL*	N3C	
+23	ARS	18	
+24	CAS	SC1	
+25	TRA	*-4	
+26	TRA	*+2	
+27	TRA	*-6	
+28	SXD	*+1,1	
+29	TXL	SKT1,4,**C	NEW ENTRY
+30	SXD	SKT2+1,1	CHANGING OLD ENTRY.
+31	SXA	SC1,4	
+32	LXA	SC1,2	
SKT2	TXI	*+1,2,-1	FIRST,
+1	TXL	SKT3,2,**0	DELETE
+2	CLA*	N3C	CLC
+3	LDQ*	N3B	ENTRY
+4	STC*	N3B	
+5	XCL		
+6	STC*	N3C	
+7	TXI	SKT2,4,-1	NEXT, DELETE
SKT3	CLA	NCDES	FRCM
+1	SUB	SC1	SINK
+2	LXA	SINK,4	TABLE
+3	SXD	SKT4,4	
+4	PAX	,6	
+5	TXI	*+1,6,1	
+6	CLA*	SK1C	
+7	SUB*	SK2C	
+8	STC	SC1	
+9	SUB	SKCUM	
+10	SLW	SKCUM	
SKT5	TXI	*+1,2,1	

```

SKT4   TXF      SKT6,2,**C
      +1 CLA*    SK1B
      +2 SUB     SC1
      +3 STC*    SK1C
      +4 CLA*    SK2B
      +5 SUB     SC1
      +6 STC*    SK2C
      +7 TXI     SKT5,4,1
SKT6   ZET      SC2
      +1 TRA     SKT7
      +2 TXI     **1,4,-1
      +3 SXA     SINK,4
      +4 TXI     **1,1,1
      +5 SXA     SKH,1
      RETURN    SKTLU
      +6 TRA     SKTLU+1
SKT1   CLA*    N3C
      +1 LDQ*    N3A
      +2 STC*    N3A
      +3 XCL
      +4 STD*    N3C
      +5 TXI     **1,1,-1
      +6 SXA     SKH,1
      +7 LXA     SINK,4
      +8 TXI     **1,4,1
      +9 SXA     SINK,4
SKT7   CLA     SKCUM
      +1 STO*    SK2C
      +2 CLA     SC2
      +3 TOV     **1
      +4 ACC     SKCUM
      +5 TNO     **2
      +6 TSX     ERROR,4
      +7 STO     SKCUM
      +8 STO*    SK1C
      RETURN    SKTLU
      +9 TRA     SKTLU+1

```

CHANGING ENTRY.
DELETING ENTRY.

SCTLU PERFORMS SAME FUNCTION AS SKTLU FR SOURCES

```

SCTLU  BEGIN    3,7
      +8 TXL     **5,0,0
      +9 CLA     1,4
      +10 STO    SC1
      +11 TNZ    **2
      +12 TSX    ERROR,4
      +13 CAS    NODES
      +14 TSX    ERROR,4
      +15 NOP
      +16 CLA     2,4
      +17 TSX    FRACT,4
      +18 TSX    ERROR,4

```

SUBROUTINE LINKAGE

```

+18 STQ SC2
+19 LXA SCH,1
+20 CLA SC1
+21 PAX ,4
+22 CAL* N3C
+23 ANA K.M6
+24 CAS SC1
+25 TRA *-4
+26 TRA *+2
+27 TRA *-6
+28 SXD *+1,1
+29 TXL SCT1,4,**0
+30 SXC SCT2+1,1
+31 SXA SC1,4
+32 LXA SC1,2
SCT2 TXI *+1,2,-1
+1 TXL SCT3,2,**0
+2 CLA* N3C
+3 LDQ* N3B
+4 STA* N3B
+5 XCL
+6 STA* N3C
+7 TXI SCT2,4,-1
SCT3 CLA NODES
+1 SUB SC1
+2 LXA SOURCE,4
+3 SXD SCT4,4
+4 PAX ,6
+5 TXI *+1,6,1
+6 CLA* SC1C
+7 SUB* SC2C
+8 STO SC1
+9 SUB SCCUM
+10 SLW SCCUM
SCT5 TXI *+1,2,1
SCT4 TXH SCT6,2,**C
+1 CLA* SC1B
+2 SUB SC1
+3 STO* SC1C
+4 CLA* SC2B
+5 SUB SC1
+6 STO* SC2C
+7 TXI SCT5,4,1
SCT6 ZET SC2
+1 TRA SCT7
+2 TXI *+1,4,-1
+3 SXA SOURCE,4
+4 TXI *+1,1,1
+5 SXA SCH,1
RETURN SCTLU
+6 TRA SCTLU+1
SCT1 CLA* N3C
+1 LDQ* N3A
+2 STA* N3A

```

NEW ENTRY
CHANGING OLD ENTRY.

FIRST,
DELETE
OLD
ENTRY

NEXT, DELETE
FRM
SCLRCE
TABLE

CHANGING ENTRY.
DELETING ENTRY.

```
+3 XCL
+4 STA* N3C
+5 TXI **1,1,-1
+6 SXA SCH,1
+7 LXA SOURCE,4
+8 TXI **1,4,1
+9 SXA SOURCE,4
SCT7 CLA SCCUM
+1 STO* SC2C
+2 CLA SC2
+3 TOV **1
+4 ADD SCCUM
+5 TNO **2
+6 TSX ERROR,4
+7 STO SCCUM
+8 STO* SCIC
RETURN SCTLU
+9 TRA SCTLU+1
```

SIDIST PRINTS DISTRIBUTION OF PATH-LENGTHS
OF MSGS DELIVERED BY SIMULATOR.

```

      XEJECT
SIDIST STL  SYSOED
      XPRINT J,1
      +3 STL  SYSOED
      +6 PZE  F10.3,,2
      BEGIN  1,4
HOD    TXL  **+3,0,0          SUBROUTINE LINKAGE
      +4 CLA  HODA
      +5 TSX  PB,4
      +6 TSX  HODX,4
      RETURN HOD
      +7 TRA  HOD+1

```

```

      BEGIN  1,7
HODX   TXL  **+5,0,0          SUBROUTINE LINKAGE
      +8 LDQ  ARRAYG
      +9 MPY  ARRAYG
      +10 LRS  10
      +11 XCA
      +12 SUB  MU
      +13 SLW  MU
      XPRINT K,1
      +14 STL  SYSOED
      +17 PZE  F11.1,,12
      XPRINT A,1
      +18 STL  SYSOED
      +21 PZE  COUNTS,,7
      RETURN HODX
      +22 TRA  HODX+1

```

FLOW DISPLAYS FLOW THRU NETWORK. MESSAGES PASSED THRU
NODES ARE DISPLAYED ROW-WISE, BEGINNING WITH FIRST ROW.

```

      BEGIN  1,7
SIFLOW TXL  **+5,0,0          SUBROUTINE LINKAGE
FLWO   EQU  SIFLOW
      -16 NZT  SNAP
      -15 TRA  1,4
      XEJECT
      -14 STL  SYSOED
      XPRINT J,1
      -11 STL  SYSOED
      -8  PZE  F10.6,,2
      -7  CLA  N4A
      -6  STA  FLW1
      -5  LXA  COLUMN,4
      -4  SXD  FLW1,4

```

```

-3  SXD  FLW2,4
-2  LXA  NODES,4
-1  TXL  FLW0+1,4,0
FLW3  CLA  FLW1
+1  SUB  COLUMN
+2  STA  FLW1
    XPRINT H,1
+3  STL  SYSOED
FLW1  PZE  **0,,**0
FLW2  TIX  FLW3,4,**0
    RETURN  FLW0
+1  TRA  FLW0+1

```

SIFLOW PRINTS DIST OF DELAYS IN STACKS FOR DLVC MSGS

```

XEJECT
SIWAIT  STL  SYSOED
    XPRINT  J,1
+3  STL  SYSOED
+6  PZE  F10.5,,2
+7  CLA  STDA

```

```

PB      BEGIN  1,7,1
    TXL  **7,0,0
+11  STA  PB1
+12  STZ  MU
+13  LXA  MAXHO,1
+14  CLA*  PB1
+15  TNZ  PB2
+16  TIX  *-2,1,1
    RETURN  PB
+17  TRA  PB+1
PB2    SXD  PB3,1
+1  AXT  1,1
PB1    LDQ  **0,1
+1  PXA
+2  XCA
+3  DVP  ARRCNT
+4  TSX  PBX,4
+5  TXI  **1,1,1
PB3    TXL  PB1,1,**0
    RETURN  PB
+1  TRA  PB+1

```

SUBROUTINE LINKAGE

```

PBX    BEGIN  1,7
    TXL  **5,0,0
+8  STQ  DT+1
+9  MPY  K.A100
+10 RND
+11 PAX  ,2

```

SUBROUTINE LINKAGE

```

+12 TXH    **2,2,0
+13 TXI    **1,2,1
+14 AXT    17,4
+15 CAL    FILL-6
+16 TIX    **3,2,6
+17 CAL    FILL,2
+18 AXT    ,2
+19 SLW    OT.1+18,4
+20 TIX    *-4,4,1
+21 SXA    OT,1
      XPRINT C,1
+22 STL    SYSOFC
+25 PZE    OT,,2C
+26 LDQ    OT
+27 MPY    OT
+28 MPY    OT+1
+29 LLS    10
+30 ADD    MU
+31 STO    MU
+32 STL    PTTGL
      RETURN PBX
+33 TRA    PBX+1

```

```

SETHMX SETS HMAX TO VALLE WHICH INSURES LESS THAN
A FRACTION, K, OF DRCP-CUTS.
CALL SEQU    TSX SETHMX,4
      PZE LOC OF K,TAG
      RETURN IF NOGC
      NORMAL RETURN

```

```

      BEGIN    3,7,1
SETHMX TXL    **7,0,0
      CLA*    1,4
+12 XCL
+13 PXA
+14 LLS    8
+15 SUB    KINT2
+16 STA    **4
+17 TZE    **2
+18 TPL    2,4
+19 PXA
+20 LRS    **0
+21 STQ    K.T1
+22 STZ    K.T2
+23 LXA    MAXHD,4
SET1 LDQ*    HODC
+1 PXA
+2 XCA
+3 DVP    ARRCNT
+4 XCA

```

SUBROUTINE LINKAGE

```

+5 ADD     K.T2
+6 STC     K.T2
+7 CAS     K.T1
+8 TXI     SET2,4,1
+9 TXI     SET2,4,1
+10 TIX    SET1,4,1
SET2 TXH    *+3,4,63
+1 SXA     HMAX,4
    RETURN  SETHMX
+2 TRA     SETHMX+1
    RETURN  SETHMX,1
+3 AXT     1,4

```

```

CLEAR ZEROES ALL COUNTS
CALL SEQU    TSX CLEAR,4
RETURN

```

```

CLEAR BEGIN 1,4
+4 TXL     *+3,0,0
+5 LXA     MAXHO,4
+6 STZ*    HODC
+7 STZ*    STDC
+7 TIX     *-2,4,1
+8 LXA     NODES,1
+9 STZ*    N2A
+10 ZET    SNAP
+11 STZ*   N4A
+12 TIX    *-3,1,1
+13 STZ    DRPCNT
+14 STZ    STKCNT
+15 STZ    ARRCNT
+16 STZ    ARRHO1
+17 STZ    ARRHO
    RETURN  CLEAR
+18 TRA    CLEAR+1

```

SUBROUTINE LINKAGE

BESTHO PRESETS NOTABLE TO BEST-PATHS, THEN
INVOKES MODIST(BELOW)

```
CALL SEQU TSX  BESTHC,4
          PZE  L,TAG
          RETURN IF NET CAN NOT BE RUN
          NORMAL RETURN
WHERE (L,TAG) CONTAINS
      LOCATION OF FIRST SPECIAL LINKS,,NR OF SPECIAL LINKS
AND L=TAG=0 IMPLIES NO SPECIAL LINKS.
```

```
BESTHO BEGIN 3,7,1
BESTHO TXL  **7,0,0          SUBROUTINE LINKAGE
+11 STL  ACTTGL          NO ACTIVITY
+12 CLA  1,4
+13 TZE  **2
+14 CLA* 1,4
+15 STO  **2
+16 TSX  PRESET,4        INIT.NET
+17 PZE
+18 TRA  **3            NO,GO
+19 TSX  MODIST,4
      RETURN  BESTHO
+20 TRA  BESTHO+1
      RETURN  BESTHO,1
+21 AXT  1,4
```

MODIST COMPUTES AND DISPLAYS THE BEST-PATH
DIST. FOR THE NOTABLE EXTANT.

```
CALL SEQU TSX  MODIST,4
          RETURN
```

```
MODIST BEGIN 1,7,1
MODIST TXL  **7,0,0          SUBROUTINE LINKAGE
HTD EQU  MODIST
-11 TSX  HTF,4            COMPUTE FOD(I)
-10 TSX  HTS,4            COMPUTE HO STAT.
      XEJECT
-9 STL  SYSOED
      XPRINT J,1
-6 STL  SYSOED
-3 PZE  F10.4,,2
-2 TSX  MOD,4            PRINT
      RETURN  HTD
-1 TRA  HTD+1
```

	BEGIN	1,7	COMP. HO-DIST.
HTF	TXL	**5,0,0	SUBROUTINE LINKAGE
+8	TSX	CLEAR,4	
+9	LXA	NODES,1	
HTO	SXA	HT1,1	
+1	PXA	,1	
+2	ALS	3	
+3	PAX	,1	
+4	SXD	HT2,1	
+5	TXI	**1,1,7	
HT3	CLA*	LINEA	
+1	TMI	HT2	
+2	STA	HT4	
+3	SXA	HT5,1	
+4	AXT	,1	
+5	LXA	NODES,2	
+6	AXT	4,4	
HT4	LDQ	**0,1	
+1	PXA		
+2	LGL	9	
+3	NZT*	N2B	
+4	TRA	**4	
+5	CAS*	N2B	
+6	TRA	**3	
+7	TRA	**2	
+8	STO*	N2B	
+9	TNX	**3,2,1	
+10	TIX	HT4+1,4,1	
+11	TXI	HT4-1,1,-1	
HT5	AXT	**0,1	
HT2	TXL	**2,1,**0	
+1	TXI	HT3,1,-1	
+2	LXA	NODES,1	
HT7	CLA*	N2A	
+1	ARS	3	
+2	CAS	MAXHO	
+3	CLA	MAXHO	
+4	NOP		
+5	PAX	,2	
+6	CLA*	HODB	
+7	ADD	K.A1	
+8	STO*	HODB	
HT6	STZ*	N2A	
+1	TIX	HT7,1,1	
HT1	AXT	,1	
+1	TIX	HTO,1,1	
	RETURN	HTF	
+2	TRA	HTF+1	

```
          BEGIN      1,7      COMPUTE AVG AND COLNT CF HC-TABLE
HTS      TXL        *+5,0,0      SUBRCUTINE LINKAGE
      +8  LXA        MAXHO,3
      +9  PXA
      +10 ADD*      HODA
      +11 TIX        *-1,1,1
      +12 STO        ARRCNT
      +13 STZ        ARRAVG
HT8      PXA        ,2
      +1  XCA
      +2 MPY*      HODB
      +3  XCA
      +4  ADD        ARRAVG
      +5  STO        ARRAVG
      +6  TIX        HT8,2,1
      +7  XCA
      +8  PXA
      +9  LLS        10
      +10 DVP       ARRCNT
      +11 STQ       ARRAVG
          RETURN    HTS
      +12 TRA       HTS+1
```

TS COMPUTES AND PRINTS M2'S DIST AS
FUNCTION OF IMPARE

CALL SEQU TSX MODEL2,4
PZE L,TAG
RETURN IF NET CAN NOT BE RUN
NORMAL RETURN

WHERE (L,TAG) CONTAINS
LOCATION OF FIRST SPECIAL LINKS,,NR OF SPECIAL LINKS
AND L=TAG=0 IMPLIES NO SPECIAL LINKS.

	BEGIN	3,7,1	
MODEL2	TXL	*+7,0,0	SUBROUTINE LINKAGE
+11	STL	ACTTGL	NO ACTIVITY
+12	CLA	1,4	
+13	TZE	*+2	
+14	CLA*	1,4	
+15	STC	*+2	
+16	TSX	PRESET,4	INIT. NET
+17	PZE		
+18	TRA	TSX	NO GO
TS	EQU	MODEL2	
	XEJECT		
-41	STL	SYSOED	
	XPRINT	J,1	
-38	STL	SYSOED	
-35	PZE	F10.2,,2	
	XPRINT	1,1	
-34	STL	SYSOED	
-31	PZE	IMPARE,,1	
-30	CLA	IMPARE	
-29	TSX	FRACT,4	
-28	TRA	TSX	
-27	STQ	K.PI1	
-26	CLA	KFULL	
-25	SUB	K.PI1	
-24	STC	K.PI	
-23	TSX	HTF,4	
-22	TSX	HTS,4	
-21	LXA	MAXHO,1	
-20	NZT*	HODA	
-19	TIX	*-1,1,1	
-18	SXA	TS1,1	
-17	LDQ*	HODA	
-16	PXA		
-15	XCA		
-14	DVP	ARRCNT	
-13	STQ*	HODA	
-12	STZ*	STDA	
-11	TIX	*-6,1,1	
-10	AXT	,1	
-9	STZ*	STDA	

```

-8 STZ MU
-7 STZ ARRCNT
-6 STZ ARRAVG
-5 AXT 1,1
-4 CLA K.PI
-3 STC* STDA
-2 PXA ,1
-1 STO TS.T1
TS1 AXT **0,1
+1 STZ K.T2
+2 LDQ* HCDA
+3 MPY* STDA
+4 ADD K.T2
+5 STO K.T2
+6 TIX *-4,1,1
+7 TZE TS3
+8 LXA TS.T1,1
+9 STO TS.T2
+10 XCA
+11 TSX PBX,4
+12 PXA ,1
+13 XCA
+14 MPY TS.T2
+15 LLS 10
+16 ADD ARRAVG
+17 STO ARRAVG
+18 TXI **1,1,1
+19 SXA TS.T1,1
+20 LXA TS1,3
TS2 TXI **1,2,-1
+1 LDQ* STDA
+2 MPY K.PI1
+3 STO TS.T2
+4 LDQ* STDB
+5 MPY K.PI
+6 ADD TS.T2
+7 STO* STDA
+8 TIX TS2,1,1
+9 TRA TS1
TS3 TSX HODX,4
RETURN TS
+1 TRA TS+1
RETURN TS,1
TSX AXT 1,4

```

MD COMPUTES AND PRINTS M1'S DIST AS
AS FUNCTION OF IMPARE

CALL SEQU TSX MODEL1,4
PZE L,TAG
RETURN IF NET CAN NOT BE RUN
NORMAL RETURN

WHERE (L,TAG) CONTAINS
LOCATION OF FIRST SPECIAL LINKS,,NR OF SPECIAL LINKS
AND L=TAG=0 IMPLIES NO SPECIAL LINKS.

	BEGIN	3,7,1	
MODEL1	TXL	**7,0,0	SLBRCLTINE LINKAGE
+11	STL	ACTTGL	NO ACTIVITY
+12	CLA	1,4	
+13	TZE	**2	
+14	CLA*	1,4	
+15	STO	**2	
+16	TSX	PRESET,4	INIT. NET
+17	PZE		
+18	TRA	MDX	NO GO
MD	EQU	MODEL1	
-13	CLA	IMPARE	
-12	TSX	FRACT,4	
-11	TRA	MDX	
-10	STQ	MD.0	
-9	MPY	MD.0	
-8	STC	MD.1	
-7	XCA		
-6	MPY	MD.0	
-5	STO	MD.2	
-4	TSX	HTF,4	
-3	TSX	HTS,4	
-2	AXT	1,1	
-1	STZ*	STDA	
MDO	LDC*	HODA	
+1	PXA		
+2	XCA		
+3	TZE	MD1	
+4	STZ*	HODA	
+5	DVP	ARRCNT	
+6	XCA		
+7	ADC*	STDA	
+8	PBT		
+9	TXI	**2,1,1	
+10	TRA	MD2	
+11	STO*	STDA	
+12	TRA	MDO	
MD1	TXI	MD2,1,-1	
MD2	SXA	MD3,1	
+1	TXI	**1,1,-1	
+2	SXD	MD6,1	

```

+3 LXA MAXHO,1
+4 SXD MD8,1
+5 SXD MD5+1,1
MD9 AXT 32000,4
+1 LDQ RND
MD3 AXT **0,1
+1 MPY K.R
+2 TRA **2
MD4 TXI **+1,1,-1
+1 CLA* STDA
+2 TLQ MD4
+3 AXT ,2
MD5 TXI **+1,2,1
+1 TXH MD8+1,2,**0
+2 MPY K.R
+3 CLA MD.0
+4 TLQ **+2
+5 TRA MD7
+6 CLA MD.1
+7 TLQ **+2
+8 TXI MD5+1,2,1
+9 CLA MD.2
+10 TLQ MD5+2
MD6 TXH MD5,1,**C
+1 TXI MD5,1,1
MD7 TIX MD5,1,1
MD8 TXL **+2,2,**0
+1 LXA MAXHO,2
+2 CLA* HOOB
+3 ADD K.A1
+4 STO* HOOB
+5 TIX MD3,4,1
+6 STQ RND
+7 TSX HTS,4
XEJECT
+8 STL SYSOED
XPRINT J,1
+11 STL SYSOED
+14 PZE F10.1,,2
XPRINT I,1
+15 STL SYSOED
+18 PZE IMPARE,,1
+19 TSX HOD,4
RETURN MD
+20 TRA MD+1
RETURN MD,1
MDX AXT 1,4

```

```

FRACT CONVERTS FLOATING FRACTION TO BINARY CNE
CALL SEQU TSX FRACR,4 (ACC) = FLTNG NR
RETURN IF GREATER THAN 1

```

NORMAL RETURN

(MQ) = BINARY NR

FRACT	XCL	
+1	PXA	
+2	LLS	8
+3	SUB	KINT2
+4	TZE	2,4
+5	TPL	1,4
+6	STA	**2
+7	PXA	
+8	LRS	**0
+9	TRA	2,4

MAKE POSITIVE.

DEFINE FORMATS

	BEGIN	1,4
FORMAT	TXL	**3,0,0
	XFORM	11
+4	STL	SYSOED
+7	PZE	FM1,,4
+8	PZE	FM2,,3
+9	PZE	FM3,,4
+10	PZE	FM4,,4
+11	PZE	FM5,,3
+12	PZE	FM6,,3
+13	PZE	FM7,,3
+14	PZE	FM8,,3
+15	PZE	FM9,,5
+16	PZE	FM10,,5
+17	PZE	FM11,,2
	RETURN	FORMAT
+18	TRA	FORMAT+1

SUBROUTINE LINKAGE

PCOL

LIMIT	EQU	32767	
TPU	EQU	20	
PCOL		*,,.POOL-*	
K.A1		1	
K.A101	PZE	1,,1	
K.A2		2	
K.A3		3	
K.A7		7	
K.A8		8	
K.A9	DEC	9	
K.A16	DEC	16	
K.A100	DEC	100	
K.A48	DEC	48	
KINT1	OCT	2330000000000	
KINT2	OCT	200	
KINT3	OCT	1000000000	
KFULL	OCT	377777777777	
K.MZE	MZE		
K.PI	DEC	0	
K.P11	DEC	0	
K.R	OCT	11060471625	MLLT FOR RANDOM
K.T1			
K.T2			
K.T3			
K.T4	PZE		
K.M1	OCT	77700000	
K.M2	OCT	77777777	
K.M3		,7	TAG
K.M4	OCT	77700000	
K.M5	OCT	777	
K.M6	OCT	7777	
K.M7	EQU	K.M5	
K.D			DESTINATION
K.HO			HO NUMBER
K.L			LINE NR.
K.MSG			MESSAGE NR.
K.N			NODE NR.
K.O			ORIGIN
K.P			PRIORITY
K.ST			TIME DELAYED IN STACKS
K.AL			COUNTER
ACTTGL	PZE	0	
	BSS	8	
AL	OCT	0	
ALHO	OCT	0	
ALWD	OCT	0	
AR			ARRIVAL LIST
BH.A	PZE	0	
BH.M	PZE	0	
B.HO	PZE	0	
BFTST	PZE	0	
BOTTCM		**0,2	

B.T.	OCT	1600004C6C00	
CONN	EQU	WEAVE	
CX	EQU	COLUMN	
DELIVR	TSX	DAO,4	
DSPY	PZE	0	
D1	OCT	0	
DR			CRCP LIST
+1	BCI	1,*****	
+2	BCI	1,*****	
+3	BCI	1,****	
+4	BCI	1,***	
+5	BCI	1,**	
+6	BCI	1,*	
FILL	BCI	1,	
FACTOR	DEC	0	
FM1	BCI	4,A,2,10X,4I11,3(F10.4B25)	
FM3	BCI	4,C,1X,I10,F12.9B0,A1C8	
FM2	BCI	3,B,5X,1C(I10)	
FM4	BCI	4,D,40X,A6,3H = ,F9.7	
FM5	BCI	3,E,40X,A6,3H = ,15	
FM6	BCI	3,F,40X,A6,3H = ,112	
FM7	BCI	3,G,40X,A6,3H = ,A6	
FM8	BCI	3,H,2,10X,12I10	
FM9	BCI	5,I,40X,9HIMPARE = ,F7.4	
FM10	BCI	5,J,40X,A12,13H DISTRIBUTION	
F10.1	BCI	2,MODEL-M1	
F10.2	BCI	2,MODEL-M2	
F10.3	BCI	2,SIMULATOR	
F10.4	BCI	2,BEST-PATH	
F10.5	BCI	2,STACK-DELAY	
F10.6	BCI	2,TRAFFIC-FLOW	
FM11	BCI	2,K,2,14X,A72	
F11.1	BCI	9,TIME HO-DROPS STACK-DROPS DELIVERIES PC-DROPPED DI	
+9	BCI	1,ST-AVG	
+10	BCI	2, VARIANCE	
YES	BCI	1, YES	
NO	BCI	1, NO	
HA.1			
MOHIGH	OCT	776	
MOH11	OCT	770	
MOINC		8	
MOINC1	EQU	MOINC	
STKINC	OCT	1000	
MOLINE			
HOMAX			
HOMEL			
HOVER	TSX	HA0,4	
+1	NOP		
+2	ARS	9	
+3	ARS	18	
HSH	ARS	27	
+1	NOP		
+2	LGL	9	
+3	LGL	18	

MSH1	LGL	27
+1	ALS	15
+2	ALS	6
+3	ARS	3
MSH2	ARS	12
IHA		
IHB		
IHC		
IHD	DEC	0
IL.T	DEC	0
INITGL		
K1	DEC	0
K2	DEC	0
LRN1	DEC	0
LEFT		**0,2
LOWER	PZE	MSG2
	VFD	15/0,3/2,15/0,3/1
+1	VFD	15/0,3/4,15/0,3/3
+2	VFD	15/0,3/6,15/0,3/5
+3	VFD	15/1,3/0,15/3,3/7
.LK1		
+1	OCT	0,0,0
.LK2	OCT	10,0,-10,0
.LD	OCT	0,0,0,0
.LS	OCT	0
MAXACT	DEC	0
MD.0		
MD.1		
MD.2		
MOVER	TSX	MA0,4
NG3	DEC	0
NG4	DEC	0
OT	DEC	0
+1	DEC	.2580
CT.1	BCI	,
+10	BCI	9,
OL	PZE	0
PRETGL	OCT	0
PTTGL	PZE	0
RA.0	OCT	0
RA.1	OCT	0
RA.2	OCT	0
RA.3	OCT	0
RA.4	OCT	0
RA.5	OCT	0
RND		
RX	EQU	ROW
ROUTER	TSX	RA0,4
SC1		
SC2		
SCCUM	OCT	0

SCH
 SIDES OCT 34000700CCO
 SINK OCT 0
 SKCUM OCT 0
 SKH
 SOURCE OCT 0
 SSLIM DEC 10
 TPHIGH PZE TPU
 TS.T1 DEC 0
 TS.T2 DEC 0
 UPPER PZE LIMIT
 UPDATE TLQ HA1
 COUNTS
 DRPCNT
 STKCNT
 ARRCNT
 CROPPC
 ARRAVG
 MU PZE 0
 HOAVG
 ARRHCI
 ARRHCI
 FONCNT
 MOTOT1
 MOTOT
 CNTTBL PZE COUNTS,,*-COUNTS
 .POOL *-POOL

NR MSGS DRCPED EN-RCUTE
 NR. MSGS DRCPED IN STACK
 NR OF ARRIVALS
 PERCENT DRCPED
 AVG. HO CF DELIVERED MSGS
 SECOND MCMENT OF DIST
 AVG. OF HO TABLE
 TOTAL HO OF DELIVERED MSGS
 NR. OF ENTRIES IN HO-TABLE
 SUM OF ENTRIES IN HO-TABLE

PARAMETER LIST

```

PARAM      *...PARAM--1
CHOKE      PZE      *
SNAP       PZE      *
FLOW       EQU      SNAP
BIND       PZE      0
ADAPT      PZE      0
ALPHA      DEC      0
IMPARE     DEC      0
IMPAIR     EQU      IMPARE
LEARN      DEC      1.0
FORGET     DEC      0
HMAX       DEC      63
PAXHC      EQU      HMAX
WEAVE      DEC      0
ROW        DEC      0
COLUMN     DEC      0
TPMAX      DEC      6
TPMIN      DEC      0
STACK      DEC      8
HPRIME     DEC      0
INITPO     EQU      HPRIME
GRAIN      OCT      2
RANDCM     OCT      110604716255
LINKS      DEC      0
.PARAM

```

TABLES

```

BEGINNING OF TABLES AND LISTS
*...MSG2--*+1

```

LINK TABLE - CONTAINS IMAGE LINK NRS. AND
DISTANCE FROM EFF TO AFF NODE
LINKS ARE NUMBERED 0 THRU 7 , VIZ.

7 0 1
6 * 2
5 4 3

LET L BE A LINE NR., THEN
 $L' = \text{EFF. NODE NR.} = \text{GRTST. INTEGER } (L/8)$
 $L'' = \text{EFF. LINK NR.} = \text{REMAINDER } (L/8)$
 $L1' = \text{AFF. NODE NR.} = L' + \text{ADDRESS PART OF LINK}(L'')$
 $L1'' = \text{IMAGE LINK NR.} = \text{FIRST 3 BITS OF LINK}(L'')$

LINK	FOR	LINK TABLE	
		**0	-D
+1	FVE	**0	1-D
+2	SIX	**0	1
+3	SVN	**0	D+1
+4	PZE	**0	D
+5	PON	**0	-(1-D)
+6	PTW	**0	-1
+7	PTH	**0	-(D+1)

TP LIST CONTAINS POINTERS TO FIRST AND LAST
MESSAGES ON TRAFFIC LISTS. TP(I') IS HEAD OF LIST
FOR TIME I' , WHERE
I' IS COUNTED MODULO TPU, THE MAXIMUM PIPE-FILLING TIME.
TP(I') = PZE LAST MSG NR , , FIRST MSG NR
TP = T' , WHERE T IS CURRENT TIME.

TPTBLE	BSS	TPU,0
TP	CCT	0

CONNECTIVITY TABLE
 CTABLE HAS ENTRIES CORRESPONDING TO ALLOWABLE
 NETWORK CONNECTIVITIES (SEE LINK TABLE)

CONN	LINKS USED	CTABLE(CONN)
1	6,2	10111C11
2	6,4,2,C	1C1C1C1C
3	7,6,4,3,2,C	CC1CCC1C
4	ALL	0C0C0C0C

ETC., AS REQUIRED

+1	OCT	31400000C000	CONN = 7
+2	OCT	146000000000	CONN = 6
+3	OCT	-20000000C000	CONN = 5
+4	OCT	0	CONN = 4
+5	OCT	10400000C000	CONN = 3
+6	OCT	-124000000000	CONN = 2
+7	OCT	-166000CCCC00	CONN = 1
CTABLE	OCT	-377000000000	ERROR
NOW , NUMBER OF LINES PER NODE FOR GIVEN CONN.			
+1	DEC	2	CONN = 1
+2	DEC	4	CONN = 2
+3	DEC	6	ETC.
+4	DEC	8	
+5	DEC	6	
+6	DEC	4	
+7	DEC	4	

LINE TABLE
PREFIX = BUSY SIGNAL
DECR = PIPE-FILLING TIME, TP
TAG = ZERO
ADDR = POINTER TO FIRST HCTABLE ENTRY FOR THIS LINE
(ZERO IF NO ENTRIES)

LINES	DEC	0
LINE1		**0,2
LINEA		**0,1
LINEB		**0,2
LINEC		**0,4
LINTBL		**0,,**0

NODE TABLE
THREE PARTS

PART ONE

PREFIX = 4 IF KILLED, ZERO IF ALIVE
DECREMENT = STANDARD-STACK POINTER
TAG = ZERO
ADDRESS = NEW-MESSAGE-STACK POINTER

PART TWO

CURRENT DEPTH OF STANDARD STACK

PART THREE

ADDRESS (DECREMENT) CONTAINS REFERENCE NODE NUMBER
FOR CHOOSING ORIGINS (DESTINATIONS).

PART FOUR EXISTS ONLY IF FLW SNAPSHOTS ARE BEING TAKEN,
AND THEN CONTAINS NR. OF MESSAGES PASSED THRU NODE.

NODES	DEC	0
N1A		**0,1
N1B		**0,2
N1C		**0,4
N2A		**0,1
N2B		**0,2
N2C		**0,4
N3A	PZE	**0,1
N3B	PZE	**0,2
N3C	PZE	**0,4
N4A	PZE	**0,1
N4B	PZE	**0,2
N4C	PZE	**0,4
NODTBL		**0,,**0

BUSY TABLE

PREFIX = PZE IMPLIES LINE IS DEAD
DECR = TIME REMAINING BEFORE LINE IS FREE
ADDR = PCINTER TO IMAGE LINE

BUSY	PZE	**0,,**0
BUSYA	PZE	**0,1
BUSYB	PZE	**0,2
BUSYC	PZE	**0,4

HANDOVER NUMBER TABLE
 INITIALIZED BY ROUTINES INH AND BNF
 PACKED FOUR/WORD, ORDERED, (FOR A GIVEN LINE)
 BY INCREASING NODE NR AND READ FORWARD (SCUTH)

HO TBL **0,,**0

HO DISTRIBUTION TABLE - READ BACKWARD
 LINE I OF HO TBL CONTAINS COUNT OF DELIVERED
 MSGS. WHOSE HO NRS. LIE BETWEEN I AND I+1

HODA **0,1
 HOCB **0,2
 HOC **0,4
 HO TBL **0,,**0

STD TBL CONTAINS DISTRIBUTION OF DELAYS IN STACKS FOR CLVD MSGS

STDA **0,1
 STCB **0,2
 STC **0,4
 STD TBL **0,,**0

SOURCE/SINK TABLE CONTAINS PERCENTAGE OF
ORIGINS/DESTINATIONS TO BE PREEMPTED BY
SOURCE/SINK NODES. UPPER PERC. BOUND
IN PART 1, LOWER IN, PART 2.

SCTBLE	**0,0,**0
SC1A	**0,1
SC1B	**0,2
SC1C	**0,4
SC2A	**0,1
SC2B	**0,2
SC2C	**0,4
SKTBLE	**0,0,**0
SK1A	**0,1
SK1B	**0,2
SK1C	**0,4
SK2A	**0,1
SK2B	**0,2
SK2C	**0,4

MESSAGE TABLE -- TWO PARTS

MSG1 CONTAINS

7 BITS - ORIGIN NODE NR.
 7 BITS - DESTINATION NODE NR.
 4 BITS - PRIORITY
 9 BITS - TIME DELAYED IN STACKS
 9 BITS - CURRENT FC

MSG2 CONTAINS

PREFIX = ZERO IF MOVING AND ALIVE, FOUR IF STACKED
 DECREMENT = POINTER TO NEXT MESSAGE ON TRAFFIC LIST
 FOR SAME TIME-FRAME, ZERO IF LAST. SEE TP DESCRIPTION
 TAG = CURRENT LINK NR.
 ADDRESS = CURRENT LINE NR.

LOAD	DEC	0
MSGCNT	DEC	0
MSG5	DEC	0
BASE		MSG2+1
MSGTBL		MSG2,,**C
MSG1		
MSG2		
END	NETPGM	

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ACTIVE 0049	SOURCE C076	CH2 *0047	FM7 C074	IAL 0021
ACTIPL 0073	STDTBL C084	CH3 C047	FM8 C074	IAL0 0024
ARRAVG 0076	STKNT C076	CHAIN C047	FM9 C074	IAL3 0024
ARRCNT 0076	STKNC C074	CHCKE 0077	FRACT C074	IAL4 0024
ARRH01 0076	SYSBFD C000	CLEAR 0064	GM C050	IAL5 0024
BESTHO 0065	SYSD81 0000	COAN 0074	GM1 C050	IAL6 0024
BOTTOM 0073	SYSD82 C000	COAT 0053	GM2 C050	IAL7 0024
CNTTBL 0076	SYSDSK C000	COX 0074	GM3 C050	IAL8 0024
COLUMN 0077	SYSLBC C000	D1 0074	GM4 C050	IAL9 0025
COUNTS 0076	SYSMIT C000	DA *0038	GM5 C050	IA20 0025
CTABLE 0080	SYSMCT C000	DAC 0038	GM6 C051	IA21 0025
DELIVR 0074	SYSD8D C000	DAL *0038	GM7 C051	IA3 0021
DROPPC 0076	SYSDRG C000	DA2 0038	GM8 C050	IA4 0021
DRPCNT 0076	SYSSBF C000	DA3 0038	GM9 C051	IA5 0021
FACTOR 0074	SYSTEM C000	DA4 0038	GN C046	IA6 0021
FORGET 0077	SYSTST C000	DR C074	GRAIN C077	IA7 0021
FORMAT 0072	TABLES C077	DRC C044	HA *0035	IA8 0021
H0D1ST 0065	TPHIGH C076	DRCP C044	HAO C035	IA9 0021
H0D1BL 0084	TPTBLE *C079	DSPY 0074	HAI C035	IA9 0021
H0HIGH 0074	UPDATE C076	DUMP 0074	HA2 C035	IH C026
H0INCL 0074	.ENTRY *C010	ERRDR C007	HA3 C035	IH1 C026
H0LINE 0074	.ERRCR *C01C	F1C.1 C074	HA4 C035	IH2 C026
H0NCNT 0076	.PARAM C077	F1C.2 C074	HA5 C035	IH3 C026
H0TBL 0084	.ACTVI C049	F1C.3 C074	HA.1 C074	IH4 C027
H0TOT1 0076	ADAPT 0077	F1C.4 C074	HMAX C077	IH5 C027
HPRIME 0077	AL C073	F1C.5 C074	HCAVG C076	IH6 0027
IMPAIR *0077	ALMC C073	F1C.6 C074	HCD C061	IH7 C028
IMPARE 0077	ALPHA 0077	F11.1 C074	HCD8 0084	IH9 0027
INITGL 0075	AR C073	FA1 C042	HCD8 0084	IHA 0075
INITHO 0077	ARRHO C076	FA2 C042	HCD8 0084	IHR C075
K-A100 0073	BASE 0086	FA3 C042	HCD8 0084	IHC C075
K-A101 0073	BASE 0086	FA6 C042	HCD8 0084	IHD 0075
LINTBL 0081	BFTST 0073	FB C043	HCD8 0084	IHX 0028
MAXACT 0075	BH C029	FB1 C043	HCD8 0084	IL C023
MODEL1 0070	BH1 C029	FB2 C043	HCD8 0084	IL0 0023
MODEL2 0068	BH10 C030	FB3 C043	HCD8 0084	IL00 0023
MSGCNT 0086	BH2 C029	FB4 C043	HCD8 0084	IL1 0023
MSGTBL 0086	BH3 C029	FB5 C043	HCD8 0084	IL2 0023
NETPGM 0007	BH4 C029	FB6 C043	HCD8 0084	IL3 0023
NODTBL 0082	BH5 C029	FILL 0074	HCD8 0084	IL4 0023
PARAMS 0011	BH6 C029	FLC *0077	HCD8 0084	ILX 0024
PARSET *0008	BH8 C030	FLCW 0061	HCD8 0084	IL.T C075
PRESET 0011	BH9 C030	FLW0 0061	HCD8 0084	INL 0016
PRETGL 0075	BHX C030	FLW1 C062	HCD8 0084	INL0 0016
RAND00 0077	BH.A C073	FLW2 C062	HCD8 0084	INL1 0016
RECORD 0045	BH.P C073	FLW3 C062	HCD8 0084	INL2 0016
ROUTER 0075	BINC C077	FMI C074	HCD8 0084	INL3 0016
SIDIST 0061	BUSY C083	FMI0 0074	HCD8 0084	INS 0017
SIFLOW 0061	BUSYA C083	FMI1 0074	HCD8 0084	INS0 0017
S1WAIT 0062	BUSYB C083	FMI2 0074	HCD8 0084	INSX 0020
SCTBLE 0085	BUSYC C083	FM3 0074	HCD8 0084	K1 C075
SETHMX 0063	B.HC 0073	FM4 0074	HCD8 0084	K2 0075
SIMUL8 0053	B.F. C074	FM5 0074	HCD8 0084	KFULL 0073
SKTBL 0085	CH1 C047	FM6 0074	HCD8 0084	KINTL 0073

KINT2	0073	LNNO	0036	NG2	0055	RA6	*C04C	SK2B	0085
KINT3	0073	LRN1	0075	NG3	0075	RA7	C04C	SK2C	0085
K.A1	0073	MA	*C032	NG4	0075	RA8	C040	SKCUM	0076
K.A16	0073	MA0	C032	NG5	0055	RA9	C040	SKH	0076
K.A2	*0073	MA1	C032	NGA	C034	RAX	C039	SKT1	0058
K.A3	0073	MA2	C033	NGB	0054	RA.0	C075	SKT2	C057
K.A4B	0073	MA2.1	C033	NGC	0053	RA.1	0075	SKT3	C057
K.A7	0073	MA3	C032	NGD	C054	RA.2	*0075	SKT4	0058
K.A8	0073	MA4	C032	NO	C074	RA.3	*C075	SKT5	C057
K.A9	0073	MAXHO	C077	NCCES	C082	RA.4	C075	SKT6	0058
K.AL	0073	MD	0070	CL	0075	RA.5	0075	SKTLU	0057
K.D	0073	MCO	0070	OT	0075	RCO	C045	SNAP	0077
K.HO	0073	MCL	C070	OT.1	0075	RM	C052	SSLIM	0076
K.L	0073	MD2	C070	PARAM	0077	RM1	C052	STACK	0077
K.M1	0073	MD3	0071	PB	0062	RM2	C052	STDA	0084
K.M2	0073	MD4	0071	PB1	0062	RM3	C052	STDB	0084
K.M3	0073	MD5	C071	PB2	0062	RND	C075	STDC	0084
K.M4	*0073	MD6	C071	PB3	0062	RCM	C077	TP	0079
K.M5	0073	MD7	0071	PBX	0062	RX	0075	TPMAX	0077
K.M6	0073	MD8	0071	PMC	0015	SC1	C075	TPMIN	0077
K.M7	0073	MD9	*C071	PMC	0014	SC1A	C085	TPU	0073
K.MSG	0073	MDX	C071	PCCL	0073	SC1B	C085	TS	0068
K.MZE	0073	MD.0	0075	PRO	C011	SC1C	0085	TS1	0069
K.N	0073	MD.1	0075	PR1.1	0012	SC2	C075	TS2	0069
K.O	0073	MD.2	C075	PR4	0014	SC2A	C085	TS3	0069
K.P	0073	MG	C048	PR5	0013	SC2B	0085	TSX	0069
K.PI	0073	MG1	0048	PR6	0012	SC2C	0085	TS.T1	0076
K.PI1	0073	MIN	0036	PR7	0012	SCCLM	CC75	TS.T2	0076
K.R	0073	MOVER	0075	PR8	0011	SCH	CC76	UPPER	0076
K.ST	0073	MSG1	C086	PRES	0011	SCT1	0059	WEAVE	0077
K.T1	0073	MSG2	0086	PR51	C0C8	SCT2	0059	YES	*0074
K.T2	0073	MSG3	0086	PR52	00C8	SCT3	C059	.FNL	C01C
K.T3	0073	MSG5	0086	PR53	00C8	SCT4	C059	.EN2	0010
K.T4	0073	MU	C076	PR54	0009	SCT5	0059	.EN3	0010
LEARN	0077	N1A	C082	PR5X	0009	SCT6	C059	.EN4	001C
LEFT	0075	N1B	0082	PRX	0014	SCT7	C060	.EN5	C01C
LIMIT	0073	N1C	C082	PTTGL	0075	SCTLU	C058	.ENX	C010
LINE1	0081	N2A	C082	RA	*0039	SET1	C063	.EXIT	0010
LINEA	0081	N2B	0082	RA0	0039	SET2	C064	.LD	0075
LINEB	0081	N2C	C082	RAC0	C039	SIDES	C076	.LK1	0075
LINEC	0081	N3A	C082	RA01	0039	SIM1	C053	.LK2	0075
LINES	0081	N3B	C082	RA02	0039	SIM2	0053	.LS	0075
LINES	0081	N3C	C082	RA03	0039	SINK	CC76	.PCCL	0076
LINKS	0077	N4A	0082	RA1	0039	SK1A	CC85		
LINKS	0077	N4B	0082	RA2	*0039	SK1B	C085		
LOAD	0086	N4C	C082	RA3	0040	SK1C	C085		
LOWER	0075	NG0	0053	RA4	0040	SK2A	C085		
LRN	0036	NG1	0056						

ON DISTRIBUTED COMMUNICATIONS:

List of Publications in the Series

- I. Introduction to Distributed Communications Networks, Paul Baran, RM-3420-PR.
Introduces the system concept and outlines the requirements for and design considerations of the distributed digital data communications network. Considers especially the use of redundancy as a means of withstanding heavy enemy attacks. A general understanding of the proposal may be obtained by reading this volume and Vol. XI.
- II. Digital Simulation of Hot-Potato Routing in a Broadband Distributed Communications Network, Sharla P. Boehm and Paul Baran, RM-3103-PR.
Describes a computer simulation of the message routing scheme proposed. The basic routing doctrine permitted a network to suffer a large number of breaks, then reconstitute itself by rapidly relearning to make best use of the surviving links.
- III. Determination of Path-Lengths in a Distributed Network, J. W. Smith, RM-3578-PR.
Continues model simulation reported in Vol. II. The program was rewritten in a more powerful computer language allowing examination of larger networks. Modification of the routing doctrine by intermittently reducing the input data rate of local traffic reduced to a low level the number of message blocks taking excessively long paths. The level was so low that a deterministic equation was required in lieu of Monte Carlo to examine the now rare event of a long message block path. The results of both the simulation and the equation agreed in the area of overlapping validity.

IV. Priority, Precedence, and Overload, Paul Baran, RM-3638-PR.

The creation of dynamic or flexible priority and precedence structures within a communication system handling a mixture of traffic with different data rate, urgency, and importance levels is discussed. The goal chosen is optimum utilization of the communications resource within a seriously degraded and overloaded network.

V. History, Alternative Approaches, and Comparisons, Paul Baran, RM-3097-PR.

A background paper acknowledging the efforts of people in many fields working toward the development of large communications systems where system reliability and survivability are mandatory. A consideration of terminology is designed to acquaint the reader with the diverse, sometimes conflicting, definitions used. The evolution of the distributed network is traced, and a number of earlier hardware proposals are outlined.

VI. Mini-Cost Microwave, Paul Baran, RM-3762-PR.

The technical feasibility of constructing an extremely low-cost, all-digital, X- or K_u -band microwave relay system, operating at a multi-megabit per second data rate, is examined. The use of newly developed varactor multipliers permits the design of a miniature, all-solid-state microwave repeater powered by a thermoelectric converter burning L-P fuel.

VII. Tentative Engineering Specifications and Preliminary Design for a High-Data-Rate Distributed Network Switching Node, Paul Baran, RM-3763-PR.

High-speed, or "hot-potato," store-and-forward message block relaying forms the heart of the proposed information transmission system. The Switching Nodes are the units in which the complex processing takes place. The node is described in sufficient engineering detail to estimate the components required. Timing calculations, together with a projected implementation

scheme, provide a strong foundation for the belief that the construction and use of the node is practical.

VIII. The Multiplexing Station, Paul Baran, RM-3764-PR.

A description of the Multiplexing Stations which connect subscribers to the Switching Nodes. The presentation is in engineering detail, demonstrating how the network will simultaneously process traffic from up to 1024 separate users sending a mixture of start-stop teletypewriter, digital voice, and other synchronous signals at various rates.

IX. Security, Secrecy, and Tamper-Free Considerations, Paul Baran, RM-3765-PR.

Considers the security aspects of a system of the type proposed, in which secrecy is of paramount importance. Describes the safeguards to be built into the network, and evaluates the premise that the existence of "spies" within the supposedly secure system must be anticipated. Security provisions are based on the belief that protection is best obtained by raising the "price" of espied information to a level which becomes excessive. The treatment of the subject is itself unclassified.

X. Cost Estimate, Paul Baran, RM-3766-PR.

A detailed cost estimate for the entire proposed system, based on an arbitrary network configuration of 400 Switching Nodes, servicing 100,000 simultaneous users via 200 Multiplexing Stations. Assuming a usable life of ten years, all costs, including operating costs, are estimated at about \$60,000,000 per year.

XI. Summary Overview, Paul Baran, RM-3767-PR.

Summarizes the system proposal, highlighting the more important features. Considers the particular advantages of the distributed network, and comments on disadvantages. An outline is given of the manner in which future research aimed at an actual implementation of the network might be conducted. Together with the introductory volume, it provides a general description of the entire system concept.